# AFFINE DUAL EQUIVALENCE AND k-SCHUR FUNCTIONS

### SAMI ASSAF AND SARA BILLEY

ABSTRACT. The k-Schur functions were first introduced by Lapointe, Lascoux and Morse [20] in the hopes of refining the expansion of Macdonald polynomials into Schur functions. Recently, an alternative definition for k-Schur functions was given by Lam, Lapointe, Morse, and Shimozono [19] as the weighted generating function of starred strong tableaux which correspond with labeled saturated chains in the Bruhat order on the affine symmetric group modulo the symmetric group. This definition has been shown to correspond to the Schubert basis for the affine Grassmannian of type A [17] and at t=1 it is equivalent to the k-tableaux characterization of Lapointe and Morse [24]. In this paper, we extend Haiman's dual equivalence relation on standard Young tableaux [14] to all starred strong tableaux. The elementary equivalence relations can be interpreted as labeled edges in a graph which share many of the properties of Assaf's dual equivalence graphs. These graphs display much of the complexity of working with k-Schur functions and the interval structure on  $S_n/S_n$ . We introduce the notions of flattening and squashing skew starred strong tableaux in analogy with jeu da taquin slides in order to give a method to find all isomorphism types for affine dual equivalence graphs of rank 4. Finally, we make connections between k-Schur functions and both LLT and Macdonald polynomials by comparing the graphs for these functions.

### 1. Introduction

Classically, the Schur functions have played a central role in the theory of symmetric functions [28]. They also appear in geometry as representatives for Schubert classes in the cohomology rings of Grassmannian manifolds, and they appear in representation theory as the Frobenius characteristics of irreducible  $S_n$  representations and as the trace for certain irreducible  $GL_n$  representations.

In [20], Lapointe, Lascoux and Morse introduced a new larger family of symmetric functions which includes the Schur functions, namely the k-Schur functions, with similar connections both to geometry and to representation theory. The k-Schur functions were defined in hopes of refining and ultimately proving the Macdonald Positivity Conjecture [27]. Precisely, Lapointe Lascoux and Morse conjectured that the Macdonald polynomials expand into k-Schur functions with polynomial coefficients in two parameters q, t with nonnegative integer coefficients, and that the k-Schur functions expand into Schur functions with polynomial coefficients with parameter

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t and nonnegative integer coefficients. Haiman [13] has since shown that the Macdonald polynomials are the Frobenius characteristic of a bigraded  $S_n$ -module defined by Garsia and Haiman [7] using the geometry of the Hilbert scheme of points in the plane. This resolved the n! Conjecture and provided the first proof of Macdonald positivity.

At this time, a number of conjecturally equivalent definitions for k-Schur functions exist [19, 20, 21, 22, 23, 24], making the term "k-Schur function" rather ambiguous. In this paper, we advocate for the geometrically inspired definition as the weighted generating function of starred strong tableaux presented by Lam, Lapointe, Morse and Shimozono [19]. This definition at t=1 is equivalent to the k-tableaux characterization in [24] which has been shown to represent the Schubert basis in the homology of the affine Grassmannian of type A [17]. Furthermore, the starred strong tableaux are a natural generalization of standard tableaux which appear throughout combinatorics.

Recently, Lam, Lapoint, Morse and Shimozono proved that the k-Schur functions as defined below except with t = 1 are Schur positive [18]. Their approach shows how k-Schur functions relate to k + 1-Schur functions when the t is not included.

It is an open problem to show that the k-Schur functions including the t statistic are Schur positive. Toward proving this conjecture, we define a family of involutions on starred strong tableaux which generalize Haiman's elementary dual equivalence moves on standard Young tableaux [14]. Using these involutions, one can put a graph structure on starred strong tableaux which satisfies many of the same axioms as the dual equivalence graphs defined by the first author in [1]. As our model for dual equivalence is based on the poset on n-cores induced from Young's lattice, our results extend to k-Schur functions indexed by skew shapes. Our main result is that these graphs, which we call affine dual equivalence graphs, are locally Schur positive when restricted to edges of 2 adjacent colors and the spin is constant on connected components, see Definition 4.5 and Theorem 7.15.  $^{1}$ 

Jeu da taquin is an important algorithm in the theory of symmetric functions related to Littlewood-Richardson coefficients. One of the properties of jeu da taquin slides is that they commutes with elementary dual equivalence moves on tableaux [14, Lemma 2.3]. There is no known analog of jeu da taquin for k-Schur functions at this time. Such an analog would in principle be useful for multiplying k-Schur functions and expanding again into k-Schurs. One approach to finding such a jeu da taquin algorithm is to look for sliding moves which commute with affine dual equivalence moves. In Sections 7.1 and 7.2, we describe two types of collapsing moves which commute with affine dual equivalence in specified cases. These collapsing moves are the analogs of removing empty rows and columns in a skew tableau via jeu da taquin.

<sup>&</sup>lt;sup>1</sup>Earlier, we announced the stronger result that k-Schur functions as defined here are Schur positive. However, we have since realized that the proof is incomplete for two reasons. First, the proof outline requires one to identify all isomorphism types for 3-colored components in affine dual equivalence graphs of the form. Our computer verification relies on a halting problem which has not terminated. Second, the axiom (4') required in [1] is not known to hold for affine dual equivalence graphs.

One of the main consequences of our results is a connection between k-Schur functions and LLT polynomials which is realized by an isomorphism of graphs for the two functions in certain cases. More generally, we expect that a better understanding of the connections between the graph we construct for k-Schur functions and that for LLT polynomials will ultimately show that an LLT polynomial expands into k-Schur functions with coefficients that are polynomials in t with nonnegative integer coefficients for an appropriate value of k. Given Haglund's formula expanding Macdonald polynomials positively into certain LLT polynomials [9, 10], this would also establish the missing connection between Macdonald polynomials and k-Schur functions.

The outline of the paper goes as follows. In Section 2, we review the basic vocabulary on partitions, the affine symmetric group, symmetric functions and quasisymmetric functions. In particular, we review an interesting order preserving bijection between a quotient of the affine symmetric group with the n-core partitions relating Bruhat order to a subposet of Young's lattice. In Section 3, one definition of k-Schur functions expanded into fundamental quasisymmetric functions is given following [19, Conjecture 9.11. These functions can be indexed by n-cores, minimal length coset representatives for  $\widetilde{\mathcal{S}}_n/\mathcal{S}_n$ , or k=n-1 bounded partitions since all three sets are in bijection. In Section 4, we review dual equivalence on standard Young tableaux along with the associated graph structures and axioms. In Section 5, we carefully study the covering relations and the rank two intervals in the poset on n-core partitions. In Section 6, we define the affine analog of dual equivalence operations and prove these maps are involutions. The main theorem is proved at the end of Section 7. Here we also give our definition of the affine dual equivalence graph on starred strong tableaux of a given shape. In Section 8, we describe the connections between k-Schur functions and both the LLT polynomials and Macdonald polynomials. We encourage the reader to look ahead to this section after seeing the definition of k-Schur functions in Section 3 in order to see the similarities. Finally, in the Appendix, we have included some examples of k-Schur functions expanded both in quasisymmetric functions and Schur functions along with their affine dual equivalence graphs.

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### 2. Basic definitions and notations

2.1. **Partitions.** A partition  $\lambda$  is a weakly decreasing sequence of non-negative integers

$$\lambda = (\lambda_1, \lambda_2, \dots, \lambda_l), \quad \lambda_1 \ge \lambda_2 \ge \dots \ge \lambda_l > 0.$$

The Young diagram of a partition  $\lambda$  is the set of points (i, j) in  $\mathbb{N} \times \mathbb{N}$  such that  $1 \leq i \leq \lambda_j$ . We draw the diagram so that each point (i, j) is represented by the unit cell southwest of the point. Abusing notation, we will write  $\lambda$  for both the partition

and its diagram. For example, the diagram of (4, 3, 1) is



We may also represent  $\lambda$  by an infinite binary string as follows. Consider the diagram of  $\lambda$  lying in the  $\mathbb{N} \times \mathbb{N}$  plane with infinite positive axes. Walk in unit steps along the boundary of  $\lambda$ , writing 1 for each vertical step and 0 for each horizontal step. For example, (4,3,1) becomes

$$\cdots 11101010101000\cdots$$

Note that this establishes a bijective correspondence between partitions and doubly infinite binary strings s such that  $s_i = 1$  for all i < l and  $s_i = 0$  for all i > r for some  $l, r \in \mathbb{Z}$ .

For partitions  $\lambda, \mu$ , we write  $\mu \subset \lambda$  whenever the diagram of  $\mu$  is contained within the diagram of  $\lambda$ ; equivalently  $\mu_i \leq \lambda_i$  for all i. Young's lattice is defined by the partial ordering on partitions given by containment.

A standard Young tableau of shape  $\lambda$  is a saturated chain in Young's lattice from the empty partition to  $\lambda$ . As moving from rank i-1 to rank i adds a single box, filling this added box with the letter i uniquely records the chosen chain. Therefore standard Young tableaux are also characterized as bijective fillings of the cells of  $\lambda$  with the letters 1 to m so that entries increase along rows and up columns. Let  $SYT(\lambda)$  denote the set of all standard Young tableaux of shape  $\lambda$ , and let SYT denote the union of all  $SYT(\lambda)$ . For example, a standard tableau of shape (4,3,1) is

When  $\mu \subset \lambda$ , we may define the *skew diagram*  $\lambda/\mu$  to be the set theoretic difference  $\lambda - \mu$ . A standard tableau of skew shape  $\lambda/\mu$  is a saturated chain in Young's lattice from  $\mu$  to  $\lambda$ , or, equivalently, a bijective filling of the cells of  $\lambda/\mu$  with entries 1 to m so that entries increase along rows and up columns.

An addable cell for a partition  $\lambda$  is any cell c such that  $c \cup \lambda$  is again a Young diagram of a partition. Similarly, a removable cell for a partition  $\lambda$  is any cell c such that  $\lambda - c$  is again a Young diagram of a partition.

A connected skew diagram is one where exactly one cell has no cell immediately north or west of it, and exactly one cell has no cell immediately south or east of it. Two distinct connected components can meet at one point but not along an edge of a cell. A connected skew diagram is necessarily nonempty. A ribbon is a connected skew diagram containing no  $2 \times 2$  subdiagram. We may define addable and removable ribbons of  $\lambda$  just as with cells; namely, a ribbon R is an addable (resp. removable) ribbon for a partition  $\lambda$  if  $\lambda \cup R$  (resp.  $\lambda - R$ ) is again a partition.

To each cell x of a diagram  $\lambda$  associate the *content of* x defined by c(x) = i - j where the cell x lies in row j and column i. We also consider the *residue of* x, defined as the content of x modulo n. The content and residue of ribbons are defined with

respect to the southeasternmost cell. The *head* of a ribbon is its southeasternmost cell, and the *tail* of a ribbon is its *northwesternmost* cell.

The hook length of x is the number of squares above and to the right of x in  $\lambda$  including x itself. Define the bandwidth of a partition to be the number of distinct contents occupied by its cells. Equivalently, the bandwidth of a non-skew partition is its maximum hook length.

An n-core is a partition having no removable ribbon of length n. Equivalently, no hook length of  $\lambda$  is divisible by n. Young's lattice restricted to n-cores gives another ranked partial order, but it is not a lattice. This partial order on n-cores is central to the definition of k-Schur functions and strong tableaux given in Section 3.

2.2. **Affine permutations.** Here we briefly recall the necessary vocabulary on affine permutations. For a more thorough treatment of the combinatorial aspects of Coxeter groups we recommend [5], specifically see Section 8.3 for details on the affine symmetric group. Recent developments on core partitions and connections to affine Weyl groups can be found in [4, 15].

Given n, consider the set  $\widetilde{\mathcal{S}}_n$  of all bijections  $w: \mathbb{Z} \longrightarrow \mathbb{Z}$  such that

$$w(i+n) = w(i) + n \ \forall i \in \mathbb{Z}$$
 and  $w(1) + w(2) + \dots + w(n) = \binom{n+1}{2}$ .

For example, given  $i, j \in \mathbb{Z}$  such that  $i \not\equiv j$  (all congruences should be taken modulo n throughout the paper), the affine transposition  $t_{i,j} \in \widetilde{\mathcal{S}}_n$  is the periodic bijection such that  $t_{i,j}(i+p\cdot n)=j+p\cdot n$ ,  $t_{i,j}(j+p\cdot n)=i+p\cdot n$ , and  $t_{i,j}(k)=k$  for all  $k\not\equiv i$  and  $k\not\equiv j$  and all  $p\in\mathbb{Z}$ .  $\widetilde{\mathcal{S}}_n$  is known as the affine symmetric group. It is the affine Weyl group of type  $A_{n-1}$ . As a Coxeter group,  $\widetilde{\mathcal{S}}_n$  is generated by the adjacent transpositions  $s_i=t_{i,i+1}$  for  $0\leq i< n$ . If  $w=s_{i_1}s_{i_2}\cdots s_{i_p}\in\widetilde{\mathcal{S}}_n$  and p is minimal among all such expressions for w, then  $s_{i_1}s_{i_2}\cdots s_{i_p}$  is a reduced expression for w and the length of w is p, denoted  $\ell(w)=p$ . The length function is the rank function for the Bruhat order on  $\widetilde{\mathcal{S}}_n$ . As a partial order, Bruhat order can be described as the transitive closure of the relation  $w< t_{i,j}w$  if  $\ell(w)<\ell(t_{i,j}w)$ . The symmetric group  $\mathcal{S}_n$  can be viewed as the parabolic subgroup of  $\widetilde{\mathcal{S}}_n$  generated by  $s_1,\ldots,s_{n-1}$ .

Let  $Q_n$  be the minimal length coset representatives for the quotient  $S_n/S_n$ . Bruhat order restricted to  $Q_n$  is again a partial order ranked by the length function. There is a rank preserving bijection from n-core partitions to  $Q_n$  which respects the Bruhat order. This correspondence leads to useful criteria for Bruhat order on  $Q_n$  in Theorem 2.3 and the covering relation in Proposition 5.3 and Corollary 5.4. We follow [30] for terminology on partial orders.

**Definition 2.1.** [16, 23, 29] Define the function

(2.2) 
$$\mathcal{C}: Q_n \longrightarrow n\text{-core partitions}$$

recursively as follows. Associate the empty partition with the identity in  $Q_n$ ; namely,  $\mathcal{C}(\mathrm{id}) = \emptyset$ . Say  $\mathcal{C}(w) = \lambda$  and  $\ell(s_i w) > \ell(w)$ , then  $\mathcal{C}(s_i w)$  is obtained from  $\lambda$  by adding every addable cell with residue i to  $\lambda$ .

In [16, 23],  $\mathcal{C}$  is shown to be a bijection. Denote  $\mathcal{C}^{-1}$  by

(2.3) 
$$A: n\text{-core partitions} \longrightarrow Q_n.$$

Remark 2.2. Definition 2.1 can be used as an algorithm for generating n-core partitions. The reader is encouraged to look ahead to Figure 1 to see how the 3-core partitions up to length 4 are generated.

Note, if  $\lambda$  is an *n*-core with an addable cell of residue *i*, then  $\lambda$  has no removable cells of residue *i*. Similarly, if  $\lambda$  is an *n*-core with a removable cell of residue *i*, then  $\lambda$  has no addable cells of residue *i* [23, §5].

The following beautiful theorem of Lascoux shows the power of the n-core model for  $Q_n$ .

**Theorem 2.3.** [25] Given  $v, w \in Q_n$ , let  $\mu = \mathcal{C}(v)$  and  $\lambda = \mathcal{C}(w)$  be the corresponding n-core partitions. Then  $\mu \subset \lambda$  in Young's lattice if and only if v < w in Bruhat order restricted to  $Q_n$ .

2.3. Symmetric and Quasisymmetric functions. We adopt notations for the standard bases for  $\Lambda$ , the ring of symmetric functions, from [28]. For this paper, we are primarily interested in the Schur functions  $s_{\lambda}$ , indexed by partitions. The Schur functions form an orthonormal basis for  $\Lambda$  with the Hall scalar product. The Schur functions also give the irreducible characters for representations of the general linear group as well as the Schubert basis for the cohomology of the Grassmannian [6]. The k-Schur functions have analogous interpretations for each of these viewpoints.

We will use the expansion for Schur functions in terms of Gessel's fundamental quasisymmetric functions [8] rather than in terms of monomials on an alphabet  $X = \{x_1, x_2, \ldots\}$ . The k-Schur functions will have a similar expansion, presented in Section 3.3.

**Definition 2.4.** For  $\sigma \in \{\pm 1\}^{m-1}$ , the fundamental quasisymmetric function associated to  $\sigma$ , denoted  $Q_{\sigma}$ , is given by

(2.4) 
$$Q_{\sigma}(X) = \sum_{\substack{i_1 \leq \dots \leq i_m \\ i_j = i_{j+1} \Rightarrow \sigma_j = +1}} x_{i_1} \cdots x_{i_m}.$$

To connect quasisymmetric functions with Schur functions, for T a standard tableau on  $1, \ldots, m$ , define the descent signature  $\sigma(T) \in \{\pm 1\}^{m-1}$  by

(2.5) 
$$\sigma_i(T) = \begin{cases} +1 & \text{if the content of } i \text{ is less than the content of } i+1 \\ -1 & \text{if the content of } i+1 \text{ is less than the content of } i. \end{cases}$$

Note that in a standard tableau, consecutive entries may never appear along the same diagonal so the content of the cells containing i and i + 1 are never equal. In particular,  $\sigma$  is well-defined on SYT.

**Theorem 2.5.** [8] The Schur function  $s_{\lambda}$  can be expressed in terms of quasisymmetric functions by

(2.6) 
$$s_{\lambda}(X) = \sum_{T \in SYT(\lambda)} Q_{\sigma(T)}(X).$$

By Theorem 2.5, working with quasisymmetric functions instead of monomials affords us the benefit of working with standard objects instead of semistandard objects. Furthermore, the expansion in (2.6) is independent of the size of the alphabet X which could be finite or infinite.

## 3. k-Schur functions

In this section, we recall two analogs of standard Young tableaux for the n-core poset called strong tableaux and starred strong tableaux from [19]. The spin statistic is defined on starred strong tableaux. These ingredients are combined to give the definition of k-Schur functions in terms of their expansion into fundamental quasisymmetric functions.

3.1. **Strong tableaux.** Consider the poset on n-core partitions induced from Young's lattice. A strong tableau of shape  $\lambda$  is a saturated chain

$$\emptyset \subset \lambda^{(1)} \subset \lambda^{(2)} \subset \cdots \subset \lambda^{(m)} = \lambda$$

in the *n*-core poset from the empty tableau to  $\lambda$ . We denote this chain by the filling S of  $\lambda$  where all cells of  $\lambda^{(i)}/\lambda^{(i-1)}$  contain the letter i.

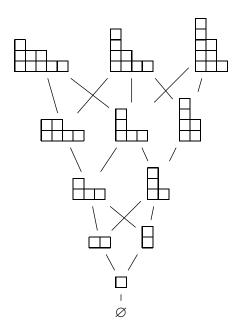


FIGURE 1. Poset of 3-cores up to rank 5.

For example, from Figure 1, the strong tableaux for n=3 of size m=4 are

						4	4
		4	4	3	3	3	3
3 4	2 4	3	2	3	2	3 4	$2 \mid 4$
1 2 3 4	1 3 3 4	1 2 3	1 3 3	$1 \mid 2 \mid 4$	1 3 4	1 2	1 3

3.2. Starred strong tableaux. A starred strong tableau,  $S^*$ , is a strong tableau  $S^*$  where one connected component of the cells containing i is chosen for each i, and the southeasternmost cell of the chosen components are adorned with a \*. Therefore, the information contained in  $S^*$  is equivalent to the pair  $(S, c^*)$  where  $c^* = (c_1, c_2, \ldots, c_m)$  is the content vector, namely  $c_i$  is the content of the cell containing  $i^*$ .

Let  $SST^*(\lambda, n)$  be the set of all starred strong tableaux of shape  $\lambda$  regarded as an n-core. For example, the 6 starred strong tableaux of shape  $\lambda = (2, 2, 1, 1)$  are

The following statistics on a starred strong tableau  $S^*$  were first introduced in [19]. Let n(i) denote the number of connected components of the cells containing i of the underlying tableau S. Among such connected components, let h(i) be the height, i.e. number of rows, of the starred connected component. Finally, let  $d(i^*)$  denote the depth of  $i^*$  in  $S^*$ , defined to be the number of components northwest of the component containing  $i^*$ . Define the statistic spin on starred strong tableaux as follows,

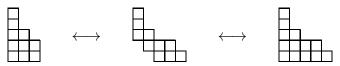
(3.2) 
$$\operatorname{spin}(S^*) = \sum_{i} n(i) \cdot (h(i) - 1) + d(i^*).$$

For example, the spins of the starred strong tableaux in equation (3.1), from left to right, are 0, 1, 1, 2, 1, 2.

This spin statistic was dubbed "spin" based on similarities with the spin statistic on ribbon tableaux that gives LLT polynomials [26]. We explore deeper connections between LLT polynomials and k-Schur functions in Section 8.

3.3. Quasisymmetric expansion. The k-Schur function  $s_{\lambda}^{(k)}(X;t)$  is the weighted generating function of starred strong tableaux of shape  $\rho(\lambda)$ , where  $\rho$  is the bijection between k-bounded partitions and k+1-cores introduced in [24]. In was also shown that the rank of  $\rho(\lambda)$  in the n-core poset equals  $|\lambda|$  and it is conjectured that the leading term of  $s_{\lambda}^{(k)}(X;t)$  in the Schur function expansion is  $s_{\lambda}(X)$ .

To define  $\rho$  on a k-bounded partition  $\lambda$ , from north to south slide each row of  $\lambda$  east as far as necessary so that no cell has hook length greater than k. Filling in the resulting skew diagram gives  $\rho(\lambda)$ . To go back, remove all cells of  $\rho(\lambda)$  with hook length greater than k and re-align the rows with the western boundary. For example, we compute  $\rho(3, 3, 2, 1, 1) = (5, 4, 2, 1, 1)$  when k = 4 as follows.



Throughout this paper, we fix n = k + 1 so that we relate n-cores with k-Schur functions.

Rather than defining a semi-standard analog of strong tableaux to define the expansion in terms of monomials as was given in [19], we formulate the definition in terms of (standard) starred strong tableaux using quasisymmetric functions. The two

versions of the definition are easily seen to be equivalent. We begin by defining the descent signature,  $\sigma \in \{\pm 1\}^{m-1}$ , of a starred strong tableau  $S^*$  of rank m as follows.

(3.3) 
$$\sigma_i(S^*) = \begin{cases} +1 & \text{if the content of } i^* \text{ is less than the content of } (i+1)^* \\ -1 & \text{if the content of } i^* \text{ is greater than the content of } (i+1)^* \end{cases}$$

Remark 3.1. Since the union of cells containing i and those containing i+1 must be a valid skew shape, the southeasternmost cells containing i and i+1 may not lie on the same diagonal. Therefore  $\sigma$  is well-defined for all starred strong tableaux.

**Definition 3.2.** Let  $\nu$  be a k-bounded partition. The k-Schur function indexed by  $\nu$  is given by

(3.4) 
$$s_{\nu}^{(k)}(X;t) = \sum_{S^* \in SST^*(\rho(\nu),n)} t^{spin(S^*)} Q_{\sigma(S^*)}(X),$$

where the sum is over all standard starred strong tableaux of shape  $\rho(\nu)$  in the n = k + 1-core poset.

Remark 3.3. We may extend Definition 3.2 to skew strong tableaux in the obvious way by considering all saturated chains from an n-core  $\mu$  to an n-core  $\nu$ . The definitions for starred strong tableaux and spin extend trivially to this setting. Consequently, all of our results for k-Schur functions also extend to this skew setting.

## 4. Dual equivalence

The main idea behind a dual equivalence graph, introduced in [2], is to provide a structure whereby the quasisymmetric functions contributing to a single Schur function are grouped together into equivalence classes, thereby demonstrating the Schur positivity of the given quasisymmetric expansion. For standard Young tableaux, the desired classes are precisely the dual equivalence classes defined by Haiman [14]. An abstract dual equivalence graph is defined by modeling the internal structure of these classes using Haiman's elementary dual equivalence relations. The connected components of a dual equivalence graph are exactly the desired equivalence classes, namely the sum over the quasisymmetric functions in a given connected component is equal to a single Schur function. Dual equivalence graphs, and more generally D graphs, provide a structure whereby we may extend the notion of dual equivalence to more general objects, in our case, starred strong tableaux.

4.1. **Dual equivalence on standard Young tableaux.** We begin by constructing a graph on standard tableaux using dual equivalence. Originally, Haiman defined an *elementary dual equivalence* on three consecutive letters i-1, i, i+1 of a permutation by switching the outer two letters whenever the middle letter is not i:

$$(4.1) \cdots i \cdots i \pm 1 \cdots i \mp 1 \cdots \cong \cdots i \mp 1 \cdots i \pm 1 \cdots i \cdots$$

In Equation (4.1),  $i \pm 1$  acts as a witness to the  $i, i \mp 1$  exchange ensuring they are not adjacent letters in the permutation.

The definition of dual equivalence extends naturally to standard Young tableaux by applying the action to the permutation obtained by reading the entries along content lines. For example, the content reading word of the standard tableau in (2.1) is 62153847. Note that in a standard tableau, i and j may lie on the same content line only if  $|i-j| \geq 3$ . In particular, each of i-1, i and i+1 must lie on distinct content lines, making equation (4.1) well-defined on standard tableaux.

It will also be helpful to think of dual equivalence on standard tableaux in terms of Young's lattice. Recall, that a standard tableau is equivalent to a saturated chain in Young's lattice with the empty partition as its minimal element. If two standard tableaux S and T are dual equivalent via an elementary dual equivalence on i-1, i, i+1, then the length two interval corresponding to the addition of i and the further away of i-1 and i+1 will be the Boolean poset on subsets of  $\{1,2\}$  ordered by containment, denoted  $B_2$ . Indeed, any length two interval in Young's lattice is either isomorphic to  $B_2$  or a chain. In this paradigm, exchanging  $i, i \neq 1$  is equivalent to traversing the length two interval where these cells are added using the other saturated chain in the interval.

We say that two standard tableaux are *dual equivalent* if one can be obtained from the other by a sequence of elementary dual equivalences. The following theorem of Haiman [14] together with Theorem 2.5 show that the sum over the quasisymmetric functions in a dual equivalence class of standard tableaux is precisely a Schur function.

**Theorem 4.1.** [14] Two standard tableaux of partition shape are dual equivalent if and only if they have the same shape.

Enrich the structure of these equivalence classes by tracking the sequence of elementary dual equivalences taking one tableau to another. Whenever T and U differ by an elementary dual equivalence for i-1, i, i+1, connect T and U with an edge colored by i. Additionally, we track the quasisymmetric function corresponding to the given tableau by writing the descent signature  $\sigma(T)$ , defined in Equation (2.5), below each tableaux. Let  $\mathcal{G}_{\lambda}$  denote the graph on all standard tableaux of shape  $\lambda$ . See Figure 2 for examples of  $\mathcal{G}_{\lambda}$ .

Define the generating function associated to  $\mathcal{G}_{\lambda}$  by

(4.2) 
$$\sum_{v \in V(\mathcal{G}_{\lambda})} Q_{\sigma(v)}(X) = s_{\lambda}(X).$$

In particular, the generating function of any vertex-signed graph whose connected components are all isomorphic to some  $\mathcal{G}_{\lambda}$  is automatically Schur positive.

4.2. **Dual equivalence graphs and D graphs.** Given any collection of objects with an associated signature function, the goal is to build a graph on the given objects that mimics the structure of these  $\mathcal{G}_{\lambda}$ . To facilitate this, we recall the local characterization of dual equivalence graphs presented in [2]. First, we need a bit of terminology.

A signed, colored graph of degree m consists of the following data: a vertex set V; a signature function  $\sigma: V \to \{\pm 1\}^{m-1}$ ; and for each 1 < i < m, a collection  $E_i$  of unordered pairs of vertices of V that represents the edges colored i. We denote such a graph by  $\mathcal{G} = (V, \sigma, E_2 \cup \cdots \cup E_{m-1})$  or simply  $(V, \sigma, E)$ .

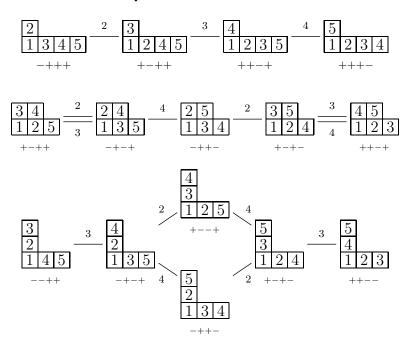


FIGURE 2. The standard dual equivalence graphs  $\mathcal{G}_{(4,1)}, \mathcal{G}_{(3,2)}$  and  $\mathcal{G}_{(3,1,1)}$ .

We say that two signed, colored graphs are *isomorphic* if there is a bijection between vertex sets that respects signatures and color-adjacency. Definition 4.2 gives criteria for when a signed, colored graph is isomorphic to  $\mathcal{G}_{\lambda}$  by Theorem 4.3.

**Definition 4.2.** A signed, colored graph  $\mathcal{G} = (V, \sigma, E)$  of degree m is a dual equivalence graph if the following hold:

- (ax1) For  $w \in V$  and 1 < i < m,  $\sigma(w)_{i-1} = -\sigma(w)_i$  if and only if there exists  $x \in V$  such that  $\{w, x\} \in E_i$ . Moreover, x is unique when it exists.
- (ax2) Whenever  $\{w, x\} \in E_i$ ,  $\sigma(w)_i = -\sigma(x)_i$  and  $\sigma(w)_h = \sigma(x)_h$  if h < i 2 or h > i + 1.
- (ax3) Whenever  $\{w, x\} \in E_i$ , if  $\sigma(w)_{i-2} = -\sigma(x)_{i-2}$ , then  $\sigma(w)_{i-2} = -\sigma(w)_{i-1}$ , and if  $\sigma(w)_{i+1} = -\sigma(x)_{i+1}$ , then  $\sigma(w)_{i+1} = -\sigma(w)_i$ .
- (ax4) For all 3 < i < m, every connected component of  $(V, \sigma, E_{i-2} \cup E_{i-1} \cup E_i)$  is either an isolated vertex or it is isomorphic to a graph in Figure 3 after the signature function is restricted to positions [i-2,i+1]. If m=4, every connected component of  $(V,\sigma,E_2 \cup E_3)$  is either an isolated vertex or it is isomorphic to a connected component in an induced subgraph of a graph in Figure 3 using only 2-edges and 3-edges and restricting the signature function to positions [i-1,i+1].

- (ax5) Whenever  $|i j| \ge 3$ ,  $\{w, x\} \in E_i$  and  $\{x, y\} \in E_j$ , there exists  $v \in V$  such that  $\{w, v\} \in E_j$  and  $\{v, y\} \in E_i$ .
- (ax6) Between any two vertices of a connected component of  $(V, \sigma, E_2 \cup \cdots \cup E_i)$ , there exists a path containing at most one edge in  $E_i$ .

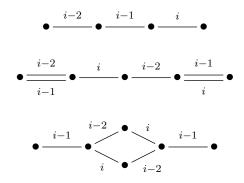


FIGURE 3. Possible 3-color connected components of a dual equivalence graph with at least two vertices. Isolated vertices are also possible.

Comparing Figure 2 with Figure 3, the largest possible connected components of  $(V, \sigma, E_{i-2} \cup E_{i-1} \cup E_i)$  are exactly the graphs for  $\mathcal{G}_{\lambda}$  when  $\lambda$  is a partition of 5. Taking this comparison to its ultimate conclusion yields the following result.

**Theorem 4.3.** [2] For  $\lambda$  a partition of m,  $\mathcal{G}_{\lambda}$  is a dual equivalence graph of degree m. Moreover, every connected component of a dual equivalence graph of degree m is isomorphic to  $\mathcal{G}_{\lambda}$  for a unique partition  $\lambda$  of m.

In practice, Axioms 1, 2 and 5 are trivially verified if  $E_i$  is the set of pairs  $\{(w, \phi_i(w)) : w \neq \phi_i(w)\}$  determined by a family of involutions  $\phi_i : V \longrightarrow V$  such that for all  $w \in V$ :

- (1) If  $\sigma(w)_{i-1,i} = +-$ , then  $\sigma(\phi_i(w))_{i-1,i} = -+$ , and vice versa.
- (2) Fixed points of  $\phi_i$  are precisely those w such that  $\sigma(w)_{i-1,i} = ++$  or --.
- (3) The signatures  $\sigma(w)$  and  $\sigma(\phi_i(w))$  agree outside the range of indices  $i-2 \le j \le i+2$ .
- (4) The involutions  $\phi_i$  and  $\phi_j$  commute whenever  $|j-i| \geq 3$ .

Axiom 3 is typically verified by keeping track of a witness in each case. The real difficulty lies in Axioms 4 and 6.

In [2], the first author extended the notion of dual equivalence in order to apply it to the LLT and Macdonald polynomials. For the extension, Axiom 6 is no longer required and Axiom 4 is replaced by a weaker axiom. Using this generalized notion of dual equivalence, it was shown that connected components have Schur positive generating functions but not necessarily a single Schur function. We use the same technique here to prove that the terms in a k-Schur function can be partitioned in connected components of a graph which are locally Schur positive. Therefore, we review the necessary material from [2] below.

**Definition 4.4.** Let  $\mathcal{G} = (V, \sigma, E)$  be a signed, colored graph. Define the generating function associated to  $\mathcal{G}$  to be

$$F_{\mathcal{G}}(X) = \sum_{v \in V} Q_{\sigma(v)}(X).$$

**Definition 4.5.** A signed, colored graph  $\mathcal{G} = (V, \sigma, E)$  of degree m is a D graph if Axioms 1, 2, 3 and 5 (from Definition 4.2) hold for  $\mathcal{G}$ . A D graph is said to be locally Schur positive on h-colored edges, denoted LSP<sub>h</sub>, provided for all  $2 \le h < i < m$ :

(LSP<sub>h</sub>) Every connected component of  $(V, \sigma, E_{i-h+1} \cup ... \cup E_i)$  using h consecutive edge sets with signatures restricted to positions [i-h, i+1] has a symmetric and Schur positive generating function.

For example, all of the graphs on Page 48 are locally Schur positive on 2-colored edges and 3-colored edges. Notice that any D graph satisfying Axioms 4 and 6 necessarily implies the graph is  $LSP_h$  for all h by Theorem 4.3.

Observe that the signature function of a D graph can be recovered from the edges plus a single sign in any one signature on any one vertex via the axioms. Thus each graph in Figure 3 can be assigned signature functions in exactly 2 ways which make them into a D graph. The third graph can only be signed in one way up to isomorphism.

#### 5. Poset on n-cores

In order to define an analog of dual equivalence for starred strong tableaux, we must first understand saturated chains in the n-core poset. In this section, we do this by exploiting the connection between n-cores and  $\widetilde{\mathcal{S}}_n$  using the abacus model for partitions.

5.1. Covering relations. We can describe the n-core poset more directly using the abacus model for cores from [16]. Consider the diagram of a partition  $\lambda$ , not necessarily an n-core, lying in the  $\mathbb{N} \times \mathbb{N}$  plane with infinite positive axes. Walk in unit steps along the boundary of  $\lambda$  placing a bead ( $\bullet$ ) on each vertical step and a spacer ( $\circ$ ) on each horizontal step. Then straighten the boundary to get a doubly infinite rod with the main diagonal marked by a vertical line. This gives the binary string uniquely representing  $\lambda$  when beads are replaced by 1's and spacers by 0's. For example, we construct the string for (4,2) as follows.

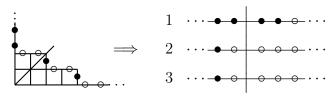


Define the *content* of a bead or spacer to be the content of the diagonal immediately southeast. Indexing each bead or spacer by its content gives an injective map from partitions to binary strings. The *abacus associated to*  $\lambda$  is the binary string of  $\lambda$  with beads and spacers indexed by their content.

Remark 5.1. Given any doubly infinite binary string s such that  $s_i$  is a bead for all i < l and  $s_i$  is a spacer for all i > r for some l, r, there is a unique re-indexing of s making it an abacus associated to a (unique) partition.

Interchanging a bead on the abacus of  $\mu$  with a spacer m places to its right corresponds to adding a ribbon of length m to  $\mu$ , and similarly interchanging a bead with a spacer m positions to its left removes an m-ribbon from  $\mu$ . In particular, if the moving bead lands in position s, then the head of the added ribbon will have content s-1.

Divide the abacus into n rods, each containing all beads and spacers of the same residue. Removing an n-ribbon from the boundary of  $\lambda$  precisely corresponds to moving a bead left along its rod. Therefore  $\lambda$  is an n-core precisely when each rod is an infinite string of beads followed by an infinite string of spacers. Define the content of a rod to be the content of the bead or spacer immediately to the right of the vertical line marking the main diagonal. We will identify a rod by its content throughout the paper. Continuing with the previous example, taking n=3 gives the following abacus decomposition of (4,2), showing rods 1,2 and 3.



Remark 5.2. Rotating the bottom row of the n-rod abacus for  $\mu$  to the top and shifting all beads in that row one column to the right will again represent the abacus for  $\mu$ , but now shifted so that the rods have contents  $0, \ldots, n-1$  from top to bottom. Similarly, rotating the top row down to the bottom and shifting all beads left on that row gives the n-rod abacus for  $\mu$  with contents  $2, \ldots, n+1$ . Thus, the abacus can be represented by n rods of contents  $k, k+1, k+2, \ldots, k+n-1$  for any integer k by scrolling up or down.

Define the *length* of each rod of the n-rod abacus as follows. For i = 1, 2, ..., n, define the length of the rod with content i to be the number of beads on the rod with positive content minus the number of spacers on the rod with nonpositive content (at most one of these numbers is nonzero). For example, the lengths of rods 1, 2, 3 for the 3-core (4, 2) are 2, -1, -1. In line with Remark 5.2, define the length of the remaining rods by setting the length of rod i - n equal to one plus the length of rod i. It is sometimes convenient to rescale the lengths of the rods so that the rods 1, 2, ..., n have nonnegative length with at least one having length 0. For now, we are concerned only with the relative lengths of the rods.

Affine permutations act on n-core partitions as discussed in Section 2.2. This action can be stated in terms of abaci as well. Recall we can represent a partition by an infinite binary string. Since affine permutations are bijections from  $\mathbb{Z}$  to  $\mathbb{Z}$ , we can apply such a bijection to any binary string. If the binary string represents an n-core then any affine transposition applied to the binary string will also represent an n-core. We leave it to the reader to verify this action is consistent with the action of simple

affine transpositions acting on n-cores described earlier. In particular, the action of an affine transposition on an n-core can be thought of pictorially as exchanging two rods of its abacus and modifying all n-translates of these two rods accordingly. The following observations, also noted in [19], follow easily from the abacus model.

**Proposition 5.3.** The following statements hold for an n-core  $\mu$  and  $t_{r,s} \in \widetilde{\mathcal{S}}_n$  with  $r < s, r \not\equiv s$ :

- (1) The abacus for  $t_{r,s}\mu$  is obtained from the abacus for  $\mu$  by swapping the lengths of the two rods with contents r and s. All rods with content distinct from r, s mod n have the same length in  $\mu$  and  $t_{r,s}\mu$ .
- (2) In the n-core poset,  $t_{r,s}\mu > \mu$  if and only if the rod of content r has larger length than the rod of content s in  $\mu$ .
- (3) An n-core  $\lambda$  covers  $\mu$  if and only if  $\lambda = t_{p,q}\mu$  for some pair p < q,  $p \not\equiv q$  such that in the abacus for  $\mu$  there is a bead at position p, a spacer at position q, and no rod between p and q has length weakly between the length of rod p and the length of rod q. Furthermore, the head and tail of one ribbon in  $\lambda/\mu$  have contents q-1 and p respectively.

*Proof.* The first statement follows form the action of an affine permutation on infinite binary strings. The second statement is immediate since moving beads right adds ribbons and moving beads left removes ribbons. The third statement also follows from this interpretation.  $\Box$ 

Proposition 5.3 is enough to describe precisely what  $\lambda/\mu$  may look like when  $\lambda$  covers  $\mu$  in the n-core poset. The condition on the lengths of the rods that lie between the interchanging rods of the abacus implies that the connected components of  $\lambda/\mu$  are identical ribbons. By Remark 5.2, the two rods being exchanged must have distinct residues and no rod between them may have the same residue as either of them. The contents across which the beads move determine the contents contained in the ribbons, and the fact that both rods are beads followed by spacers ensures that the ribbons lie on consecutive residues. These observations reprove the following statement due to Lam, Lapointe, Morse and Shimozono.

Corollary 5.4. [19, Prop. 9.5] Let  $\mu$  be an n-core and  $t_{r,s}$  an affine transposition such that  $t_{r,s}\mu$  covers  $\mu$  in the n-core poset. Then 0 < s - r < n and the connected components of  $t_{r,s}\mu/\mu$  are identical shape ribbons with cell residues from  $r \mod n$  to  $s-1 \mod n$ . Moreover, if rod r has k > 0 more beads than rod s, then  $t_{r,s}\mu/\mu$  has exactly k identical ribbons. If the head of the first ribbon lies in a cell with content c, then the head of the other ribbons have content  $c + n, c + 2n, \ldots, c + (k-1)n$ .

By Corollary 5.4, for a strong tableau S of shape  $\lambda$ , call the connected components of  $\lambda_i/\lambda_{i-1}$  the *i-ribbons of* S. Recall from Section 3.2 that a starred strong tableau consists of a strong tableau plus a choice of *i*-ribbon for each i present in S. We use the next definition and corollary to relate the starred strong tableaux to saturated chains labeled by certain sequences of transpositions.

**Definition 5.5.** Let  $\mu \subset \lambda$  be n-cores, and let  $T(\lambda/\mu, n)$  be the set of all transposition sequences

$$(t_{r_1s_1} \rightarrow t_{r_2s_2} \rightarrow \cdots \rightarrow t_{r_ms_m})$$

such that

- (1) the product  $t_{r_m s_m} \cdots t_{r_2 s_2} t_{r_1 s_1} \mu = \lambda$  as elements of  $\widetilde{\mathcal{S}}_n / \mathcal{S}_n$ ;
- (2) for each  $1 \le i \le m$ , we have  $0 < s_i r_i < n$ ;
- (3) for each  $0 \le i < m$ , the abacus for  $\mu^{(i)} = t_{r_i s_i} \cdots t_{r_2 s_2} t_{r_1 s_1} \mu$  contains a bead at position  $r_{i+1}$ , a spacer at position  $s_{i+1}$ , and every rod with content between  $r_{i+1}$  and  $s_{i+1}$  has length strictly smaller than both the length of rod  $r_{i+1}$  and the length of rod  $s_{i+1}$  or strictly larger than both.

By Proposition 5.3, condition (3) above implies  $\mu = \mu^{(0)} < \mu^{(1)} < \cdots < \mu^{(m)} = \lambda$  forms a saturated chain in the *n*-core poset. The following is a consequence of Proposition 5.3 and Corollary 5.4.

Corollary 5.6. Let  $\mu \subset \lambda$  be n-cores. There exists a bijection from skew starred strong tableaux  $S^* \in SST^*(\lambda/\mu, n)$  to  $T(\lambda/\mu, n)$  given by mapping

$$S^* \mapsto (t_{r_1s_1} \to t_{r_2s_2} \to \cdots \to t_{r_ms_m})$$

where  $s_i - 1$  and  $r_i$  are the contents of the head and tail of the i-ribbon containing  $i^*$  in  $S^*$ .

For example, this bijection maps

5.2. Intervals of length two. As motivation, recall that an elementary dual equivalence on standard tableaux may be defined in terms of interval exchanges in Young's lattice. Though the induced poset on n-cores in not as nice as Young's lattice, Björner and Brenti [5] showed that any interval of length two is either a chain or isomorphic to  $B_2$ .

**Definition 5.7.** Let  $S = (\emptyset = \mu^0 \subset \mu^1 \subset \cdots \subset \mu^m)$  be a saturated chain in the n-core poset such that the interval  $[\mu^{i-1}, \mu^{i+1}]$  is not a chain for 0 < i < m. The i-interval swap on S, denoted swap<sub>i,i+1</sub> $(S) = \text{swap}_{i+1,i}(S)$ , replaces  $\mu^i$  with the unique other n-core at rank i in  $[\mu^{i-1}, \mu^{i+1}]$ .

For example, from Figure 1 we see that a 2-interval swap on the chain

results in the chain

In terms of the strong tableaux, the same 2-interval swap gives

By Definition 5.5, the same two saturated chains can be represented by the following transposition sequence

$$\operatorname{swap}_{2,3}(t_{0,1} \to t_{1,2} \to t_{2,3} \to t_{-2,-1}) = t_{0,1} \to t_{-1,0} \to t_{1,3} \to t_{-2,-1},$$

where only the two transpositions in the middle are modified. In general, the map  $\operatorname{swap}_{i,i+1}(S)$  always modifies the *i*th and the *i*+1st transpositions in any transposition sequence representing the saturated chain S and leaves all other transpositions in the sequence fixed. The two new transpositions are not unique however since we have not yet described how the stars will move. This extension will be called a bswap and introduced in Section 6. First, we need a complete understanding of how the *i*-ribbons and the *i*+1-ribbons can appear in a strong tableau and how they change under an *i*-interval swap.

Using the abacus model for cores and Proposition 5.3, we can explicitly describe the result of an *i*-interval swap on a strong tableau  $S = (\emptyset \subset \mu^1 \subset \cdots \subset \mu^m)$  in terms of the two rod exchanges corresponding to the covering relations in the interval  $[\mu^{i-1}, \mu^{i+1}]$ . In order to analyze two consecutive rod exchanges, we extend the *n*-rod abacus picture to include extra rods above as necessary so that the four rods to be exchanged all appear as rows of the picture with the longer row above the shorter row. Ignoring the rods which are untouched by the exchange and choosing representatives of the exchanging rods as close together as possible, there are four natural cases to consider, each depicted in Figure 4: disjoint, interleaving, nested and abutting. There are three possible ways for the rod exchanges to be abutting; the two depicted and also the reverse of the right hand side. For the first three cases in Figure 4, the corresponding transpositions will have four distinct residues whereas for the abutting case, they will have only two or three distinct residues.

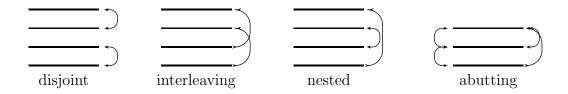


FIGURE 4. Possible rod exchanges for a length two interval.

The easiest case to consider is a disjoint exchange. Here we assume all of the residues of the rods to be exchanged are distinct, lest we actually have an abutting exchange. Further, we can assume the exchanging rods have contents a < b < c < d with n > d - a > 0 since the rods are as close together as possible. The two exchanges in this case clearly commute, and taking either first raises the rank by exactly one by Corollary 5.4. In the strong tableau, such an *i*-interval swap will happen precisely

when the cells of the *i*-ribbons and (i + 1)-ribbons have no residues in common, and the effect of the swap will be to exchange all *i*'s for i + 1's and conversely.

The case of an *interleaving exchange* is only slightly more interesting, though the conclusion of this case is noteworthy. Labeling the residues of the exchanging rods a < b < c < d from top to bottom, again we assume all four residues to be distinct lest we be pulled into the abutting case. The assumption that these two exchanges each increase the rank in the poset forces rod a longer than rod c and similarly rod b longer than rod d by Proposition 5.3. Suppose  $\mu^{i-1} \subset t_{a,c}\mu^{i-1} = \mu^i \subset t_{b,d}t_{a,c}\mu^{i-1} = \mu^{i+1}$ ; the other case is similarly resolved. By Proposition 5.3, this means the length of rod b does not lie between the lengths of rods a and c and that the length of rod a does not lie between the lengths of rods b and d. Recall, we chose a picture for the abacus so that the length of rod b is larger than the length of rod d and the length of rod a is longer than the length of rod c. These statements together imply that the lengths of rods a and c do not interleave the lengths of rods b and d, and so the transpositions taken in the other order each raise the rank by exactly one, thus  $\mu^{i-1} \subset t_{b,d}\mu^{i-1} \subset t_{a,c}t_{b,d}\mu^{i-1} = \mu^{i+1}$  is a valid strong tableau. In this new strong tableau, the contents of the *i*-ribbons and i + 1-ribbons will not overlap, though the residues will. It is also important to note that the *i*-ribbons and i+1-ribbons will not have the same residues for their heads or tails. In this case, the i-interval swap again simply exchanges all i's for i+1's and conversely. We summarize the key observation in this case as follows.

**Proposition 5.8.** An i-ribbon and an i+1-ribbon in a strong tableau have overlapping contents if and only if the contents of one ribbon are strictly contained in the contents of the other. Furthermore, the contents of the head and tail of the longer ribbon do not occur among the contents of the shorter ribbon.

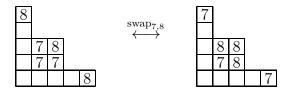
More generally, we say two ribbons are *nested* if the second condition of Proposition 5.8 holds. We also say two ribbons  $R_1$  and  $R_2$  are *independent* if  $R_1 \cup R_2$  has two connected components as a skew shape.

In the case of a *nested exchange*, again label the rod contents a < b < c < d from top to bottom. We can assume d-a < n by Corollary 5.4. Here the two corresponding transpositions commute, and each will raise the rank by exactly one. The interesting feature of this case lies in the strong tableaux.

By Proposition 5.3, neither the length of rod b nor the length of rod c may lie between the lengths of rods a and d. If both rod b and rod c are longer than rod a or both shorter than rod d (necessarily rod a is longer than rod d), then the i-ribbons and i+1-ribbons will have no contents in common, though the residues of one ribbon will be strictly contained within the residues of the other. Furthermore, both the heads and tails of the i-ribbons and i+1-ribbons have distinct residues.

On the other hand, if rod b is longer than rod a and rod c is shorter than rod d (necessarily rod b is longer than rod c), then the content of every instance of the longer ribbon (corresponding to  $t_{a,d}$ ) overlaps the content of a shorter ribbon (corresponding to  $t_{b,c}$ ) and there must be an instance of the shorter ribbon containing a cell of content b-1 which occurs independently from all of the longer ribbons. An i-interval swap changes all entries in all of the shorter ribbons that appear independently of the

longer ribbons and all entries of the longer ribbons that are not on the same content as a shorter ribbon. For example, below is the skew strong tableau corresponding to a nested exchange and the result of the interval swap



This discussion proves the following lemma.

**Lemma 5.9.** If an i-ribbon and an i+1-ribbon are nested, then

- (1) At least two copies of the shorter ribbon occur independently from the longer ribbon, with at least one on either side of the consecutive sequence of copies of the longer ribbon.
- (2) Every copy of the longer ribbon nests a copy of the shorter ribbon.
- (3) Both the heads and tails of the i-ribbons and i+1-ribbons have distinct residues.
- (4) An i-interval swap is possible.

The final case of an abutting exchange will involve exactly three distinct indices on the transpositions, though possibly only two distinct residues. Label the contents of the rods a < b < c from top to bottom. Suppose that the three residues are all distinct. This is necessarily the case for the right hand side of Figure 4. Say the two exchanges correspond with the transposition sequence  $(t_{a,c} \to t_{a,b})$ , then by Proposition 5.3, we know rod a is strictly longer than rod c which is strictly longer than rod b, and taking the inner exchange first forces rod a longer than rod b longer than rod c. Therefore we note that  $t_{b,c}t_{a,b}=t_{a,b}t_{a,c}$  so this equation along with the total order on the lengths of the rods ensures that an interval swap is possible. The new transposition sequence after applying this interval swap would be  $(t_{a,b} \to t_{b,c})$  which corresponds with the left hand side of the abutting exchange pictured in Figure 4. If the two exchanges correspond with the transposition sequence  $(t_{a,b} \to t_{a,c})$ , examining the required rod length inequalities again we see that  $(t_{b,c} \to t_{a,b})$  is a valid transposition sequence on the same rank 2 interval. This again corresponds with the left hand side of the abutting exchange pictured in Figure 4. If the right hand side of Figure 4 is turned upside down, a similar analysis holds. Furthermore, the interval swaps form an involution on the two chains in any interval isomorphic to  $B_2$  so we have covered all possible cases of an abutting exchange in the form of the left hand side of the abutting picture of Figure 4 as well. Hence in all cases of an abutting exchange with three distinct residues, there exists an interval swap determined above.

Assuming that  $[\mu^{i-1}, \mu^{i+1}]$  is isomorphic to  $B_2$ , one way to recognize if an abutting exchange is required for  $\operatorname{swap}_{i,i+1}(S)$  is that an *i*-ribbon and an i+1-ribbon together form a ribbon shape. In this case, we will say these two ribbons *abut* each other. From the transpositions pictured in the abutting case of Figure 4 and Corollary 5.4, we observe that the sum of the lengths of an *i*-ribbon and an i+1-ribbon is necessarily less than n and exactly one of the two ribbon types occurs without abutting a copy

of the other. In this instance, the *i*-interval swap will change all entries of the non-abutting ribbons and all entries in their n-translates. For example, the 2-ribbon abuts a 3-ribbon in the strong tableau on the left in (5.1).

The other way to recognize if an abutting exchange is required for  $\operatorname{swap}_{i,i+1}(S)$  is that, among the i, i+1-ribbons, one ribbon is strictly longer than the other and the longer ribbon contains an n-translate of the shorter and the heads or tails of the two ribbons have the same residue depending on if the shared rod is a or c. See for example, the 2-ribbon and 3-ribbon in the strong tableau on the right in (5.1). Here an interval swap will change all entries of the shorter ribbons and all entries of the longer ribbons that are not part of an n-translate of the shorter. This case is also recovered from the left hand side of Figure 4 when the lengths of the three rods are all distinct; we omit details as the case is completely parallel.

Following the details of the abutting exchange case carefully, we have the following.

**Proposition 5.10.** Suppose  $\operatorname{swap}_{i,i+1}(S)$  is obtained from S by an abutting exchange. Assume the corresponding transpositions are indexed by 3 distinct residues mod n. Then either

- No i-ribbon abuts any i+1-ribbon, but one of these ribbons strictly contains an n-translate of the other with a shared head or tail occurring on a consecutive residue.
- OR, all instances of one ribbon type abut the other while the other will also have at least one components which is non-abutting and the sum of the length of an i-ribbon and an i+1-ribbon is at most n-1. In this case, if an i+1-ribbon abuts an i-ribbon from the north, then the non-abutting ribbons lie always southeast of the abutting ribbons, and if an i-ribbon abuts an i+1-ribbon from the west, then the non-abutting ribbons lie always northwest of the abutting ribbons.

Finally, consider an abutting exchange as in the left hand side of the abutting case in Figure 4. If the three rod lengths are distinct and the three residues are distinct, then the exchange is covered by Prop 5.10. In each of the remaining cases, we claim the interval  $[\mu^{i-1}, \mu^{i+1}]$  is a chain so an *i*-interval swap is not possible.

**Proposition 5.11.** Let  $S = (\emptyset = \mu^0 \subset \mu^1 \subset \cdots \subset \mu^m)$  be a saturated chain in the n-core poset. Then the interval between  $\mu^{i-1}$  and  $\mu^{i+1}$  in the n-core poset is a chain if and only if each i-ribbon abuts an i+1-ribbon and each i+1-ribbon abuts an i-ribbon. Moreover, the length of an i-ribbon plus the length of a i+1-ribbon is less than or equal to n, with equality if and only if  $\mu^{i+1}/\mu^{i-1}$  is a single connected ribbon shaped component starting and ending with i+1-ribbons.

*Proof.* Assume  $[\mu^{i-1}, \mu^{i+1}]$  is a chain. Then by the previous analysis of rod exchange cases the chain corresponds with a transposition sequence of the form  $t_{b,c} \to t_{a,b}$  with a < b < c or a > b > c. Assume a and c have different residues (both necessarily have distinct residues from b). In this case,  $t_{a,b}t_{b,c} = t_{a,c}$  and by Proposition 5.3 we can assume 0 < c - a < n. Thus the skew shape  $\mu^{i+1}/\mu^{i-1}$  is the union of a positive

number of n-translates of a single ribbon of length less than n and none of these ribbons overlap in content. More precisely, i-ribbons and i + 1-ribbons always occur in pairs and the sum of their lengths is strictly less than n.

If, on the other hand, a and c have the same residue, then we can assume c = a+n by choosing to label the exchanging rods as close together as possible. Hence, the length of the ribbons corresponding to  $t_{a,b}$  and those corresponding to  $t_{b,a+n}$  necessarily add to n so  $\mu^{i+1}/\mu^{i-1}$  is a single connected ribbon shaped component. Furthermore, recall that rod c = a + n is one shorter than the length of rod a by Remark 5.2. If rod b is shorter than rod a, then the chain corresponds with the transposition sequence  $t_{a,b} \to t_{b,c}$ , otherwise the transpositions happen in the reverse order. In either case, by considering how ribbons are created using the abacus model and Proposition 5.3 we observe that the ribbon  $\mu^{i+1}/\mu^{i-1}$  is tiled by an alternating sequence of i-ribbons and i + 1-ribbons and it begins and ends with an i + 1-ribbon.

To prove the reverse direction, assume each *i*-ribbon abuts an i+1-ribbon and conversely. Then by Corollary 5.4 we can infer that the chain  $\mu^{i-1} \subset \mu^i \subset \mu^{i+1}$  corresponds to an abutting exchange. If all three contents of the exchanging rods have distinct residues, then either  $[\mu^{i-1}, \mu^{i+1}]$  is a chain or we would find a contradiction to the second case of Proposition 5.10.

If there are only two distinct indices among the exchanging rods then the relative lengths of these rods determine the only possible exchange sequence taking  $\mu^{i-1}$  to  $\mu^{i+1}$  by Proposition 5.3. Thus,  $[\mu^{i-1}, \mu^{i+1}]$  is again a chain.

Corollary 5.12. If a strong tableau  $S = (\mu^0 \subset \mu^1 \subset \mu^2)$  is the result of an abutting exchange, then  $\mu^2/\mu^0$  is the union of ribbons with nonoverlapping content. If every ribbon in  $\mu^2/\mu^0$  is an identical n-translate of the first, then the interval  $[\mu^0, \mu^2]$  is a chain.

*Proof.* This follows from the characterization of all abutting exchanges in this subsection, Proposition 5.10 and Proposition 5.11.  $\Box$ 

Table 1 summarizes the discussion above characterizing all possible length two intervals determined by two consecutive rod exchanges. Assume the initial n-core is  $\mu$ . First apply  $t_{a,b}$  then  $t_{c,d}$  assuming 0 < b - a < n, 0 < d - c < n, rod a longer than rod b, rod c longer than rod d, and all 4 indices appear in the smallest possible interval of  $\mathbb{Z}$  which satisfies these conditions. Let #res be the number of distinct residues among  $a, b, c, d \mod n$ . Let #dis be the number of distinct rod lengths among rods  $a, b, c, d \mod n$ . The interval  $[\mu, t_{cd}t_{ab}\mu]$  is either isomorphic to  $B_2$  or the chain  $C_3$  with 3 elements. The two interval types are distinguished by considering #res and #dis or equivalently by considering the skew shape as partitions of  $t_{cd}t_{ab}\mu/\mu$ .

### 6. Affine dual equivalence

We now have all the ingredients to construct an analog of dual equivalence for starred strong tableaux, which we call *affine dual equivalence*. Though our equivalence relation will not share all of the properties of dual equivalence on tableaux, we will go on in Section 7 to construct a signed colored graph from our elementary equivalence relations that we show to be a D graph.

#res	#dis	Type	Exchange	Skew Shape			
2	2	$C_3$	abutting	One long ribbon alternating $i$ 's and $i + 1$ 's,			
				starting and ending with $i+1$ .			
3	2	$C_3$	abutting	Every component is an identical ribbon			
				composed of one <i>i</i> -ribbon abutting one $i + 1$ ribbon.			
3	3	$B_2$	abutting	Long ribbons contain $n$ -translate of shorter ribbons.			
				One short ribbon occurs independently.			
	2.2.4	ъ	1	411 11 1			
4	2,3,4	$B_2$	$\operatorname{disjoint}$	All ribbons are non-overlapping,			
				i-ribbons don't abut $i + 1$ -ribbons or vice versa.			
				No $n$ -translate of one ribbon type overlaps the other.			
4	4	$B_2$	interleaving	All ribbons are non-overlapping,			
4	4	$D_2$	interieaving	i-ribbons don't abut $i + 1$ -ribbons or vice versa.			
				Some $n$ -translate of an $i$ -ribbon overlaps an			
				i+1-ribbon with distinct heads and tails.			
4	4	$B_2$	nested	Either all ribbons are non-overlapping, or			
				all longer ribbons overlap shorter ribbons and			
				at least one short ribbon occurs independently			
				NW (SE) of each long ribbon ribbons. These two			
				cases distinguished by comparing rod lengths.			

Table 1. Summary of length two interval types.

While the elementary equivalence relations will have a somewhat complicated description, there are essentially only two cases: one that precisely mirrors dual equivalence, and another that is a close approximation when the former is not applicable. Remarkably, the relations also preserve the spin statistic on starred strong tableaux.

6.1. **Elementary equivalences.** In this subsection, we describe a family of involutions  $\varphi_i$  on all starred strong tableaux of a given shape that will define the *elementary affine dual equivalence* on i-1, i, i+1. Recall that a starred strong tableau  $S^*$  of shape  $\lambda$  can be represented by a strong tableau  $S = (\emptyset \subset \lambda^{(1)} \subset \lambda^{(2)} \subset \cdots \subset \lambda^{(m)})$  with  $\lambda^{(m)} = \lambda$  and a vector  $c^* = (c_1, c_2, \ldots, c_m)$  where  $c_i$  is the content of the cell of  $S^*$  containing  $i^*$ . In this case, we will say the rank of  $S^*$  is m.

**Definition 6.1.** Let  $S^* = (S, c^*)$  be a starred strong tableau of rank m. Fix 1 < i < m. Consider the locations of  $(i-1)^*$ ,  $i^*$ ,  $(i+1)^*$  in  $S^*$ . The *i-witness*, or simply the witness when i is fixed, is chosen among  $\{i-1, i, i+1\}$  as follows.

- (1) If  $c_{i-1} \neq c_{i+1}$ , then  $c_{i-1}, c_i, c_{i+1}$  are all distinct since p-ribbons and p+1-ribbons cannot have head or tails of the same content by the analysis in Section 5.2. In this case, the *witness* is the index of the median of the set  $\{c_{i-1}, c_i, c_{i+1}\}$ .
- (2) If  $c_{i-1} = c_{i+1}$ , then we have three cases to consider.
  - (a) If the (i-1)-ribbons and (i+1)-ribbons have the same length, then i+1 is the witness.

- (b) If the (i-1)-ribbons and (i+1)-ribbons have different lengths and  $c_{i-1} > c_i$ , then the *witness* is the letter indexing the longer ribbons among the (i-1)-ribbons and the (i+1)-ribbons.
- (c) If the (i-1)-ribbons and (i+1)-ribbons have different lengths and  $c_{i-1} < c_i$ , then the *witness* is the letter indexing the shorter ribbons among the (i-1)-ribbons and the (i+1)-ribbons.

Note that when  $S^*$  is a Young tableau, the contents of the unique cells containing i-1, i and i+1 must all be distinct, ensuring that the witness is always the index of the median of the set  $\{c_{i-1}, c_i, c_{i+1}\}$ .

Next we define the involution  $\varphi_i$  on starred strong tableaux that will serve as a model for dual equivalence. Intuitively, if i and j are witnessed by h in  $S^*$ , then an elementary dual equivalence move should be based on the map  $\sup_{i,j}$  where  $\{h,j\} = \{i-1,i+1\}$ . This would be straightforward but for the difficulty of defining how the stars should behave under such an action. We obtained these rules experimentally guided by the principle that the stars should move as little as possible while preserving the spin statistic, always remaining in the same connected component of the union of cells in  $S^*$  containing i-1,i,i+1 but necessarily switching which letter they adorn. This will be the action of  $\varphi_i$  whenever such a move is possible without changing the witness. However, if the interval is a chain and the starred letters both lie in the same connected component, then neither an interval swap nor a star swap is possible. We overcome this challenge by exchanging saturated chains of length three.

**Definition 6.2.** Fix a starred strong tableau  $S^* = (S, c^*)$  of rank m with 1 < i < m. Let h be the i-witness for  $S^*$ . If  $h \neq i$ , then let j be defined by  $\{i-1, i+1\} = \{j, h\}$ . Let  $S_q$  be the union of all q-ribbons and let  $S_{q^*}$  be the connected component of  $S_q$  containing  $q^*$  for  $1 \leq q \leq m$ . We will say  $S_q$  nests  $S_{p^*}$  if the content of every cell of  $S_{p^*}$  is also the content of a cell in  $S_q$  but no head or tail of a ribbon in  $S_q$  has the same content as the head or tail of  $S_{p^*}$ . Similarly, a connected skew shape A nests another connected skew shape B provided the content of every cell of B is the content of some cell of A, but the largest and smallest contents of cells in A are not the contents of any cells in B. Let  $b_q$  be the content of the ribbon tail for  $S_{q^*}$ . We will say  $S_i$  and  $S_j$  are not abutting if  $b_i, b_j, (c_i + 1), (c_j + 1)$  have distinct residues, otherwise  $S_i$  and  $S_j$  are abutting. Let  $B_i$  and  $B_j$  be the connected components of  $S_i \cup S_j$  containing  $i^*$  and  $j^*$ , respectively.

Then  $\varphi_i(S^*)$  is defined by the first case that applies below

$$(6.1) \qquad S^* \quad \text{if } i = h,$$
 
$$\text{bswap}_{i,j}(S^*) \quad \text{if } S_i \text{ and } S_j \text{ are not abutting, or if } B_i \text{ and } B_j \text{ have different shapes and neither nests } S_{h^*},$$
 
$$\text{snake}_{i,j}^h(S^*) \quad \text{if } b_h \equiv b_j \text{ and } c_h \equiv c_j,$$
 
$$\text{bswap}_{i,j} \text{bswap}_{i,h}(S^*) \quad \text{if } S_i \text{ or } S_j \text{ nests } S_{h^*},$$
 
$$\text{double}_{i,j}^h(S^*) \quad \text{if } B_i \text{ or } B_j \text{ nests } S_{h^*},$$
 
$$\text{star}_{i,j}(S^*) \quad \text{if } B_i \neq B_j \text{ but they have the same shape.}$$
 Here the map  $\varphi_i$  depends on four types of ribbon swaps: basic swap, snake swap,

Here the map  $\varphi_i$  depends on four types of ribbon swaps: basic swap, snake swap, double swap, and star swap. Each of the ribbon swaps will only be well-defined under certain circumstances. As we prove in Theorem 6.4, the circumstances where a ribbon swap is applied in (6.1) will be precisely the circumstances when the ribbon swap is well-defined. The fact that these are all possible cases can be observed from the fact that  $h^*$  lies weakly between  $i^*$  and  $j^*$  and the notation at the beginning of Definition 6.2.

The basic swap, denoted bswap<sub>i,j</sub> $(S^*)$ , is the result of an interval swap on S and interchanging the blocks containing  $i^*$  and  $j^*$ 

$$bswap_{i,j}(S^*) = (swap_{i,j}(S), c^*(B_i \leftrightarrow B_j)).$$

For example, if 
$$n=4$$
 and  $i=4$  then  $\varphi_4=\operatorname{bswap}_{4,5}$  interchanges 
$$(6.2) \qquad \qquad \underbrace{\begin{array}{c} 5^*\\4\\1^*2^*3^*4^* \end{array}} \qquad \underbrace{\begin{array}{c} bswap_{4,5}\\4^*\\1^*2^*3^*5^* \end{array}} \qquad \underbrace{\begin{array}{c} 4\\4^*\\1^*2^*3^*5^* \end{array}}$$

In the left tableau,  $B_4$  is the cell of content 3 filled by  $4^*$  and  $B_5$  is the set of cells with contents  $\{-1, -2\}$  filled by  $4, 5^*$ . In the right tableau  $B_5$  is the cell of content 3 and  $B_4$  is the set of cells with contents  $\{-1, -2\}$ . Note, the star in the  $\{-1, -2\}$ block must move when applying the map in either direction so as to return a valid starred strong tableau with a star at the head of an i-ribbon and a j-ribbon.

A description of the operation  $c^*(B_i \leftrightarrow B_i)$  is given specifically as follows. Let  $d_p = c_p + 1$  for each p so that the p-ribbons in  $S^*$  correspond with applying the transposition  $t_{b_p,d_p}$ . Let  $r_p = d_p - b_p$  be the length of a p-ribbon in  $S^*$ . Let  $\varepsilon_p$  be the unit vector with a 1 in the p-th position. Assume p < q, then define

(6.3) 
$$\operatorname{flop}_{q,p}(c^{*}) = \operatorname{flop}_{p,q}(c^{*}) = \begin{cases} t_{p,q}(c^{*}) - r_{p} \cdot \varepsilon_{p} & \text{if } d_{p} \equiv d_{q} \text{ and } |B_{p}| < |B_{q}|, \\ t_{p,q}(c^{*}) - r_{q} \cdot \varepsilon_{q} & \text{if } d_{p} \equiv d_{q} \text{ and } |B_{p}| > |B_{q}|, \\ t_{p,q}(c^{*}) + r_{q} \cdot \varepsilon_{q} & \text{if } b_{q} \equiv d_{p} \text{ and } |B_{p}| > |B_{q}|, \\ t_{p,q}(c^{*}) + r_{p} \cdot \varepsilon_{p} & \text{if } b_{p} \equiv d_{q} \text{ and } |B_{p}| < |B_{q}|, \\ t_{p,q}(c^{*}) & \text{otherwise.} \end{cases}$$

Therefore, formally we define

$$bswap_{i,i}(S^*) = (swap_{i,i}(S), flop_{i,i}(c^*)).$$

We prove  $\operatorname{bswap}_{i,i}(S^*)$  is always a valid starred strong tableau in Theorem 6.4.

Remark 6.3. Note that when  $S^*$  is a Young tableau, it is impossible for the cell containing i to abut the cell containing j when  $h \neq i$  is the witness. Therefore the required ribbon swap in this case will always be  $\varphi_i(S^*) = \text{bswap}_{i,j}(S^*) = (\text{swap}_{i,j}(S), t_{i,j}(c^*))$ . Hence  $\varphi_i$  reduces to the usual elementary dual equivalence relation on Young tableaux.

The snake swap, denoted snake $_{i,j}^h(S^*)$ , is the result of moving the stars on all three ribbons i-1, i, i+1 while keeping the underlying strong tableau fixed. If i-1 is the witness, the moves are based on the permutation  $231 = t_{12}t_{23}$ ; if i+1 is the witness, the moves are based on the permutation  $312 = t_{23}t_{12}$ . Either way, j will become the i-witness of snake $_{i,j}^h(S^*)$ . Assuming h is the witness, then

(6.4) 
$$\operatorname{snake}_{i,j}^{h}(S^{*}) = \begin{cases} (S, \ t_{i,j}t_{i,h}(c^{*}) - r_{j} \cdot \varepsilon_{i} + r_{h} \cdot \varepsilon_{h}) & \text{if } (c_{j} < c_{i}) \text{ xor } (i < j), \\ (S, \ t_{i,j}t_{i,h}(c^{*}) + r_{i} \cdot \varepsilon_{i} - r_{i} \cdot \varepsilon_{h}) & \text{otherwise.} \end{cases}$$

We will show in the proof of Theorem 6.4 that  $\operatorname{snake}_{i,j}^h$  is only applied when  $S_i \cup S_j$  and  $S_i \cup S_h$  are both single connected ribbons so  $[\lambda^{(i-2)}, \lambda^{(i+1)}]$  is a chain by Proposition 5.11. When h = i+1, the stars move away from the diagonal of content  $c_h$  along these ribbons and when h = i-1 the stars move in toward the diagonal of content  $c_h$  along these ribbons. The star on the witness toggles between h and j by sliding along the diagonal with content  $c_h$ . For example, if n = 2 and i = 3, then  $\varphi_3 = \operatorname{snake}_{3,2}^4$  maps

4	$\operatorname{snake}_{3,2}^4$	$4^*$
3*4*	→	3 4
2 3 4	<del></del>	$2^*  3   4 $
$1^* 2^* 3 4$	$\operatorname{snake}_{3,4}^2$	$1^* 2 3^* 4$

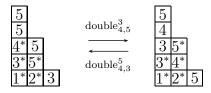
The inverse map is given by  $\operatorname{snake}_{3,4}^2$  applied to the tableau on the right.

The double swap, denoted double  $h_{i,j}(S^*)$ , is the result of two interval swaps on S and another "almost permutation" of the three relevant indices in the content vector. Precisely,

$$\operatorname{double}_{i,j}^{h}(S^{*}) = \begin{cases} \left(\operatorname{swap}_{i,j}\operatorname{swap}_{i,h}(S), \ t_{i,j}t_{i,h}(c^{*}) + r_{h} \cdot \varepsilon_{h}\right) & \text{if } b_{h} \equiv b_{j}, \\ \left(\operatorname{swap}_{i,j}\operatorname{swap}_{i,h}(S), \ t_{i,j}t_{i,h}(c^{*}) - r_{h} \cdot \varepsilon_{i}\right) & \text{if } c_{h} \equiv c_{j}. \end{cases}$$

Since double  $h_{i,j}$  is only applied when  $B_i$  or  $B_j$  nests  $S_{h^*}$  but neither  $S_i$  or  $S_j$  nests  $S_{h^*}$ , we can conclude that the nesting block is a ribbon and that  $S_j$  contains a cell with the same content as either the head or tail of  $S_{h^*}$  by considering all possible rank 2 abutting rod exchanges. Thus, when its applied either  $b_h \equiv b_j$  or  $c_h \equiv c_j$ . For example, if n = 3 and i = 4, then  $\varphi_4$  interchanges the following tableaux via double

swaps:



The star swap, denoted  $\operatorname{star}_{i,j}(S^*)$ , is the result of moving the star on  $i^*$  to the adjacent j-ribbon and vice versa while keeping the underlying strong tableau fixed. To be precise, if  $B_i$  and  $B_j$  are distinct and both  $B_i$  and  $B_j$  contain both an i and j-ribbon, then both blocks have the same shape by Proposition 5.10 and Proposition 5.11. Say f is the offset of the contents of  $B_j$  from  $B_i$ , so  $c_i + f$  is the content of the head of the i-ribbon in  $B_j$  and  $c_j - f$  is the content of the head of the j-ribbon in  $B_i$ . Then

$$\operatorname{star}_{i,j}(S^*) = (S, c^* + f \cdot \varepsilon_i - f \cdot \varepsilon_j).$$

For example, if n = 4 and i = 6, then  $\varphi_6 = \text{star}_{6,7}$  interchanges

$4^*$	$\operatorname{star}_{6,7}$	$4^*$						
3*6*77	<b>←</b>	3*	6	7	$7^*$			
1* 2* 4 5* 6 7 7*		1*	$\overline{2^*}$	4	5*	6*	7	7

6.2. A well-defined involution. Given the complicated definition of the affine dual equivalence relations, it is not obvious that  $\varphi_i$  is well-defined, much less that it is an involution. Our next task is to establish these two facts. In the course of doing so, we provide many more examples of the action of  $\varphi_i$ , though in the interest of space only the relevant cells in the strong tableaux are shown.

**Theorem 6.4.** For each 1 < i < m, the map  $\varphi_i$  is a well-defined involution on all starred strong tableaux of a fixed n-core  $\lambda$  of rank m.

*Proof.* Let  $S^*$  be a starred strong tableau of shape  $\lambda$ . We can assume  $i \neq h$  throughout the proof, the contrary case being trivial. Suppose first that  $S_i$  and  $S_j$  are not abutting. In this case, a swap<sub>i,j</sub>(S) is well defined and  $b_i, b_j, d_i, d_j$  are all distinct mod n by the classification of rod exchanges for rank 2 intervals in Section 5.2. Unless  $S_i$  and  $S_j$  come from an interleaving rod exchange with some *i*-ribbon nested in an i+1-ribbon or vice versa, the interval swap will simultaneously change all i's to j's and conversely. Therefore  $\varphi_i(S^*) = \operatorname{bswap}_{i,j}(S^*) = (\operatorname{swap}_{i,j}(S), t_{i,j}(c^*))$  is a well-defined starred strong tableau with stars in the original cells in  $S^*$ , though now adorning the opposite letter among  $\{i,j\}$  from before. When  $S_i$  and  $S_j$  come from an interleaving rod exchange with some i-ribbon nested in an i + 1-ribbon or vice versa, then the interval swap will change all entries in the shorter ribbon appearing independently as well as entries in the longer ribbon not on the same content as a shorter ribbon. In particular, the shape of the blocks  $B_i$  and  $B_j$  remains unchanged. Therefore bswap<sub>i,i</sub>( $S^*$ ) is again a valid starred strong tableau. In this case, the star adorning the longer ribbon remains in place, and the star adorning the shorter ribbon remains if the shorter ribbon is not nested in a longer, otherwise it slides one position along the diagonal; see Figure 5 for an example.

Consequently, in order to show  $\varphi_i$  is an involution in this case, it remains only to show that h remains the witness after applying  $\operatorname{bswap}_{i,j}$ . Since the effect on the content vector is merely to interchange  $c_i$  and  $c_j$ , the result follows provided  $c_h \neq c_j$ . However, the contrary case forces an i-ribbon to abut both the i-1-ribbon and i+1-ribbon with heads on content  $c_{i-1}=c_{i+1}$ . This ensures that  $\operatorname{bswap}_{i,j}$  is an involution in this case.

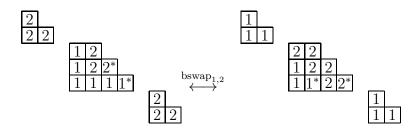


FIGURE 5. The action of  $\varphi_i$  on  $S^*$  when  $S_i$  and  $S_j$  are nested

Henceforth, we will assume that  $S_i$  and  $S_j$  are abutting, and thus both  $B_i$  and  $B_j$  must have ribbon shape by Corollary 5.12.

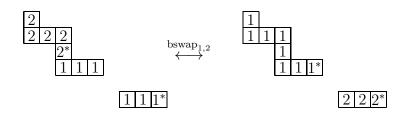


FIGURE 6. The action of  $\varphi_i$  when  $S_i \cup S_j$  is abutting and  $S_h$  is not nested.

Assume that  $B_i$  and  $B_j$  have different (ribbon) shapes but neither nests  $S_{h^*}$ . Since  $B_i$  and  $B_j$  have different shapes, the map swap<sub>i,j</sub> will toggle between each block containing only one letter and exactly one of these blocks containing both letters as shown in Figure 6. By Proposition 5.10 one can deduce how the stars move in the blocks  $B_i$  and  $B_j$  in order to adorn the other letter. These moves are summarized in the function flop<sub>i,j</sub>( $c^*$ ), hence  $T^* = \varphi_i(S^*) = \text{bswap}_{i,j}(S^*)$  is a well defined starred strong tableau. Furthermore, by inspection we have that  $\text{bswap}_{i,j}(T^*) = S^*$ . Thus,  $\varphi_i(S^*)$  is an involution provided h is also the i-witness of  $T^*$ .

Observe that the only way for the witness to change is if  $h^*$  lies on a diagonal within a block containing both i's and j's, and  $h^*$  lies weakly between their respective heads. Let I and J be the abutting i-ribbon and j-ribbon in the block overlapping  $h^*$ . By Proposition 5.8, consecutive ribbons may not have partially overlapping contents. Therefore if an h-ribbon has content overlapping an i-ribbon, one of the two must be nested. By assumption,  $S_{h^*}$  is not nested inside either  $B_i$  or  $B_j$ , hence is not nested in I. On the other hand, if h-ribbons nest i-ribbons, then they must also nest

j-ribbons, otherwise a swap<sub>i,h</sub>(S) is possible and will leave i-ribbons and j-ribbons with partially overlapping contents, again contradicting Proposition 5.8. Therefore h-ribbons and i-ribbons may not have overlapping contents, so the cell containing h\* must overlap J in content.

If J lies southeast of I, this forces the  $h^*$ -ribbon to overlap I or be nested in J, neither of which is possible. Thus J must lie northwest of I, hence the head of the h-ribbon is forced to have the same content as the head of J. Furthermore, the  $h^*$ -ribbon must be longer than J since by assumption h-ribbons are not nested in  $I \cup J$ . Therefore when  $h^*$  and  $j^*$  lie on the same diagonal, h remains the i-witness for  $\varphi_i(S^*)$  using part (2) of Definition 6.1.

Next consider the case where  $c_{i-1} \equiv c_{i+1}$  and  $b_{i-1} \equiv b_{i+1}$ . Then,  $t_{b_{i-1},d_{i-1}} = t_{b_{i+1},d_{i+1}}$  as affine permutations, where recall  $d_p = c_p + 1$ . By Corollary 5.6,  $S^*$  is associated to a transposition sequence. In order for  $t_{b_{i-1}d_{i-1}} \to t_{b_id_i} \to t_{b_{i+1}d_{i+1}}$  to be a valid triple in the transposition sequence,  $t_{b_id_i}$  must not commute with the other two. Hence at least one i+1-ribbon completely overlaps some i-1-ribbon, sharing both a head and tail, and i-ribbons must abut each such pair from both sides. By Proposition 5.11, this means  $S_i \cup S_j$  and  $S_i \cup S_h$  are both ribbons, hence snake  $b_{i,j}$  is well-defined on  $b_i$ . By inspecting (6.4), we see that  $b_i$  as its  $b_i$ -witness and snake  $b_i$  is a starred strong tableau on the same underlying strong tableau  $b_i$  with  $b_i$  as its  $b_i$ -witness and snake  $b_i$  is an involution in this case as well. For example, see Figure 7.

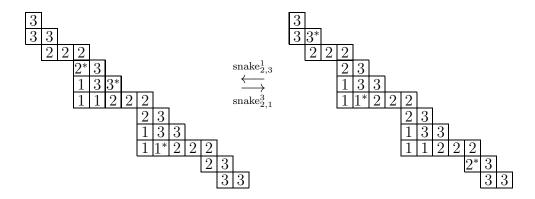


FIGURE 7. The action of  $\varphi_i$  when i-1-ribbons and i+1-ribbons share a head and tail.

Henceforth, we will assume that  $S_i$  and  $S_j$  are abutting and either  $c_{i-1} \not\equiv c_{i+1}$  or  $b_{i-1} \not\equiv b_{i+1}$ . The next case to consider is when  $B_i$  or  $B_j$  nests  $S_{h^*}$ , including the possibility that  $S_i$  or  $S_j$  nests  $S_{h^*}$ . Note that if  $B_i = B_j$ , then  $B_i$  necessarily nests  $S_{h^*}$  in order for h to be the witness.

We claim that in all these cases some connected component of  $S_h \cup S_i \cup S_j$  is a single h-ribbon. If some h-ribbon is nested in an i-ribbon or a j-ribbon, then the claim follows immediately from Lemma 5.9. So, assume that some connected component of  $S_i \cup S_j$  contains both i's and j's and nests an h-ribbon. By Proposition 5.8, we may further assume that the nested h-ribbon shares a head or tail with the j-ribbon.

Necessarily the h-ribbon and j-ribbon must both abut an i-ribbon at their shared content in the strong tableau. Therefore we have one of the following scenarios for the three transpositions corresponding to i-1, i, i+1 on the abacus, where  $a < b < c < d \le a+n$ .

To ease notation we assume h = i - 1 and j = i + 1 and we are doing the rod exchange on the left in (6.5), noting that the other cases are completely analogous. In this case, analyzing the transposition triple  $t_{b,c} \to t_{a,b} \to t_{b,d}$  shows that initially the length of  $\operatorname{rod}(b)$  cannot be weakly between the lengths of  $\operatorname{rod}(a)$  and  $\operatorname{rod}(d)$  or else the transpositions don't each increase the rank by exactly one at each step. Furthermore, the length of  $\operatorname{rod}(b)$  cannot be less than the length of  $\operatorname{rod}(d)$  because otherwise there would be no i + 1-ribbon with content overlapping any i - 1-ribbon by the definition of  $T(\lambda, n)$  and Corollary 5.6 contradicting the assumption that some connected component of  $S_i \cup S_j$  contains both i's and j's and nests an h-ribbon. Therefore, the length of  $\operatorname{rod}(b)$  is strictly greater than the length of  $\operatorname{rod}(a)$  so, by Corollary 5.6 again, there must an i - 1-ribbon occurring independently from all i, i + 1-ribbons.

With the claim proved, we conclude by Proposition 5.11 that  $T = \operatorname{swap}_{i,h}(S)$  is a well defined, valid strong tableau. After such, some *i*-ribbon appears independently of all *j*-ribbons in T, making  $U = \operatorname{swap}_{i,j}(T)$  well defined.

In the case  $B_i$  or  $B_j$  nests  $S_{h^*}$  but neither  $S_i$  or  $S_j$  nests  $S_{h^*}$ , the final result of swap<sub>i,j</sub>(swap<sub>i,h</sub>(S)) is much like a single interval swap in the case of nested i, i+1-ribbons: all independently occurring h's change to j's and all letters of  $S_i \cup S_j$  not on the same diagonal as an h will change with i's becoming h's and j's becoming i's; for example, see Figure 8. The shape and contents of the nested ribbon remains the same, but now these are j-ribbons. Therefore,  $U^* = \text{double}_{i,j}^h(S)$  is a well-defined starred strong tableau with a star placed at the head of some p-ribbon for each p. The effect of  $\varphi_i$  on the content vector for  $S^*$  is an involution by inspection. Since j-ribbons are now nested in  $U^*$ , we only need to show j becomes the i-witness in  $U^*$  in order to prove  $\varphi_i$  is an involution on such an  $S^*$ . This will clearly be the case so as long as  $c_h \neq c_j$  both before and after applying double i, Assuming i and i-ribbon will be forced to lie southeast of i0, and i1-ribbons and i1-ribbons will share a tail instead, so the witness indeed changes as desired.

Similarly, in the case  $S_i$  or  $S_j$  nests  $S_{h^*}$ , the image  $U^* = \varphi_i(S^*) = \text{bswap}_{i,j}(\text{bswap}_{i,h}(S))$  is a well-defined starred strong tableau with a star placed at the head of some p-ribbon for each p by the proof of a single basic swap above. Either  $U_i$  or  $U_h$  nests  $j^*$  and j is the i-witness for  $U^*$ . Hence  $\varphi_i^2(S^*) = S^*$  for such an  $S^*$ .

Finally, we will assume  $B_i$  and  $B_j$  have the same shape but lie on distinct content diagonals, each have ribbon shape, and neither nests  $S_{h^*}$ . Then  $\varphi_i(S^*) = \operatorname{star}_{i,j}(S^*)$ 

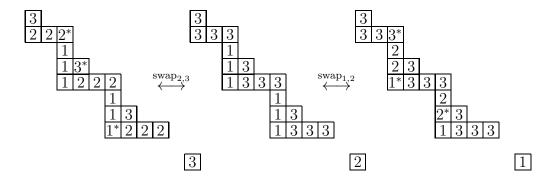


FIGURE 8. The action of  $\varphi_2 = \text{double}_{i,j}^h$  when  $S_{i^*} \cup S_{j^*}$  nests  $S_{h^*}$ .

is a well defined starred strong tableau where the same hypotheses hold. See Figure 9 for an example.

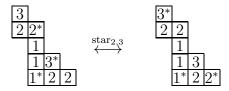


FIGURE 9. The action of  $\varphi_i$  when  $h^*$  overlaps  $S_{i^*} \cup S_{j^*}$  without being nested.

In this case, both blocks  $B_i$  and  $B_j$  contain both i's and j's. The only way for the witness to change is if  $h^*$  lies on a diagonal within a block containing both i's and j's, and  $h^*$  lies weakly between their respective heads. The proof that h remains the witness is the same as the argument above for the case  $B_i$  and  $B_j$  have different shape and neither nests  $S_{h^*}$ .

6.3. Preservation of spin. Next we show that the involution  $\varphi_i$  has the added feature of preserving the spin statistic. Recall from (3.2) that spin is defined by

$$spin(S^*) = \sum_{i} n(i) \cdot (h(i) - 1) + d(i^*),$$

where n(i) is the number of *i*-ribbons, h(i) is the height of an *i*-ribbon and  $d(i^*)$  is the depth of the starred *i*-ribbon. We will show that  $\varphi_i$  preserves the spin by tracking the contribution for i-1, i and i+1.

**Proposition 6.5.** For any starred strong tableau  $S^*$ , we have  $\text{spin}(\varphi_i(S^*)) = \text{spin}(S^*)$ .

*Proof.* Recall the notation from Definition 6.2. Assume  $i \neq h$ . If  $S_i$  and  $S_j$  are disjoint, interleaving or nested with non-overlapping content, then  $\varphi_i$  acts by simultaneously replacing all i's with j's and conversely. The contribution to spin for ribbons other

than i, j is unchanged, and these two swap contributions, thereby preserving the statistic. If any two ribbons of  $S_i$  and  $S_j$  are nested with overlapping contents, then recall that  $\varphi_i$  does not change the shape of the shorter ribbon nor the height (nor width) of the longer ribbon, and the stars remain on the same diagonals. This ensures that contributions to spin for i and j are exchanged, and all other contributions are unchanged.

We may now assume that  $S_i$  and  $S_j$  are abutting for the remainder of the proof. If  $\varphi_i$  acts by  $\operatorname{star}_{i,j}$ , then this affects only the depths of  $i^*$  and  $j^*$ . We claim  $d(i^*) + d(j^*)$  is preserved since every connected component of  $S_i \cup S_j$  inclusively between  $B_i$  and  $B_j$  have the same shape when  $\operatorname{star}_{i,j}$  is applied. Hence,  $\varphi_i$  again preserves spin.

If  $\varphi_i$  acts by a bswap<sub>i,j</sub>, then recall from the proof of Theorem 6.4 that each connected component of  $S_i \cup S_j$  is a ribbon, in particular one of  $B_i$  and  $B_j$  is a longer ribbon containing an n-translate of the shorter. Let  $n_l$  and  $n_s$  denote the number of the longer ribbons and shorter ribbons in  $S_i \cup S_j$ , respectively, and let  $h_l$  and  $h_s$  denote their respective heights. Let  $d_l$  be the number of longer ribbons northwest of the starred long ribbon, and similarly let  $d_s$  denote the number of shorter ribbons northwest of the starred short ribbon.

Supposing that the connected components of  $S_i \cup S_j$  each contain a unique letter, the contributions for i and j to spin are

(6.6) 
$$\operatorname{spin}_{S^*}(i) = n_s(h_s - 1) + d_s,$$

(6.7) 
$$\operatorname{spin}_{S^*}(j) = n_l(h_l - 1) + d_l.$$

On the other hand, letting  $T^* = \operatorname{bswap}_{i,j}(S^*)$ , some connected component of  $T_i \cup T_j$  contains both i's and j's. By Proposition 5.10, this implies that in every component containing both letters, the smaller entries are south of the larger entries if and only if the shorter ribbons appear independently to the southeast. Armed with this observation, we compute that if the longer ribbon among  $B_i$  and  $B_j$  contains both i's and j's, then the contribution to spin for i and j in  $T^*$  is

(6.8) 
$$\operatorname{spin}_{T^*}(i) = (n_s + n_l)(h_s - 1) + (d_s + \varepsilon n_l),$$

(6.9) 
$$\operatorname{spin}_{T^*}(j) = n_l(h_l - h_s + (1 - \varepsilon) - 1) + d_l,$$

where  $\varepsilon$  is 1 if the shorter ribbons appear independently southeast of the longer (equivalently, the larger entry abuts the shorter from the north), and 0 if the shorter ribbons appear independently northwest of the longer (equivalently, the smaller entry abuts the larger from the west). Noting the equality between (6.6) plus (6.7) and (6.8) plus (6.9) shows spin is once again preserved. This also handles the case when  $\varphi_i$  acts by  $\mathrm{bswap}_{i,i}\mathrm{bswap}_{i,h}$ .

Consider now the case when  $\varphi_i$  acts by double i,j. For this case, we may assume, from the analysis in Theorem 6.4, that some h-ribbon and j-ribbon share a head or tail. Also from Theorem 6.4, some h-ribbon must appear independently of all i-ribbons. Furthermore, by Proposition 5.10, if i- or j-ribbons appear independently of the other, then they do so on the opposite side of abutting i- and j-ribbons from h-ribbons. Supposing first that the combined lengths of an i-ribbon plus a j-ribbon is less than n, reading from northwest to southeast or from southeast to northwest

one sees isolated h-ribbons followed by abutting i- and j-ribbons nesting h-ribbons. There are then three options for what follows: isolated j-ribbons; isolated abutting i- and h-ribbons; or no further i-, j- or h-ribbons. For example, see Figure 10. Note that, in particular,  $S^*$  has isolated j-ribbons if and only if double i-, j- and j-ribbons.

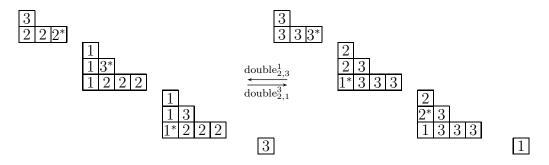


FIGURE 10. An example of  $\varphi_i$  acting via  $\varphi_i = \text{double}_{i,j}^h$ .

To assess the contributions to spin, assume that  $S^*$  has no isolated abutting i- and h-ribbons, as in the right hand side of Figure 10. Let  $h_w$  and  $d_w$  denote the height and depth of the starred h-ribbon, respectively, and let  $n_w$  denote the number of isolated witness ribbons. For the example, we have  $n_w = 1$ ,  $h_w = 1$ ,  $d_w = 0$ . Let  $n_l$  be the number of i-ribbons, each with height  $h_l$  and the starred one with depth  $d_l$ . Let  $n_s$  be the number of isolated j-ribbons, and let  $h_s$  and  $d_s$  denote the height and depth, respectively, of j-ribbons. For the example, we have  $n_l = 2$ ,  $h_l = 2$ ,  $d_l = 1$  and  $n_s = 1$ ,  $h_s = 2$ ,  $d_s = 0$ .

The contribution to spin from i-1, i, i+1 in  $S^*$ , where  $S^*$  has no isolated abutting i- and h-ribbons, is given by

(6.10) 
$$\operatorname{spin}_{S^*}(h) = (n_l + n_w)(h_w - 1) + d_w,$$

(6.11) 
$$\operatorname{spin}_{S^*}(i) = n_l(h_l - 1) + d_l,$$

(6.12) 
$$\operatorname{spin}_{S^*}(j) = (n_l + n_s)(h_s - 1) + d_s.$$

Following the description of how double i,j acts on these ribbons, we may similarly compute the contributions of i-1, i, i+1 to the spin of  $T^* = \text{double}_{i,j}^h(S^*)$ . With i and j defined relative to i, we have

(6.13) 
$$\operatorname{spin}_{T^*}(h) = (n_l + n_w + n_s)(h_w - 1) + (d_w + \varepsilon n_s),$$

(6.14) 
$$\operatorname{spin}_{T^*}(i) = (n_l + n_s)(h_s - h_w + (1 - \varepsilon) - 1) + d_s,$$

(6.15) 
$$\operatorname{spin}_{T^*}(j) = n_l(h_l + h_w - (1 - \varepsilon) - 1) + d_l,$$

where, similar to before,  $\varepsilon$  is 0 if the witness originally existed only to the left of abutting i- and j-ribbons and 1 otherwise. Adding the contributions in either case miraculously yields the same result, thereby showing that the spin statistic is preserved.

If i-ribbons and j-ribbons have lengths adding to n, we regard the abutting i-and j-ribbons which together nest an h-ribbon as abutting pairs, and the leftover  $\max(i,j)$ -ribbon as isolated. For example, in Figure 8, we regard the left side as having 1-ribbons abutting 2-ribbons from the west with an isolated abutting 2-ribbon and 3-ribbon to the northwest, and the right side we regard as having 2-ribbons abutting 3-ribbons from the west with an isolated 3-ribbon to the northwest. That is to say, Figure 8 is the same as Figure 10 for the purposes of calculating spin. In this case, note that  $S^*$  has no isolated abutting i- and h-ribbons precisely when h = i - 1. Moreover, in this case we always have  $n_s = 1$ . With this alteration, the analysis of spin is precisely as before, again showing that spin is preserved.

Finally, if  $\varphi_i$  acts by  $\operatorname{snake}_{i,j}^h$ , then the difference  $\operatorname{spin}(S^*) - \operatorname{spin}(\varphi_i(S^*))$  only depends on the change in depth for  $h^*, i^*, j^*$  since both  $S^*$  and  $\varphi_i(S^*)$  have the same underlying strong tableau. Furthermore, we have that i-1-ribbons and i+1-ribbons both have length n minus the length of an i-ribbon. In this case, there is one more i+1-ribbon than i-ribbon and one more i-ribbon than i-1-ribbon. Using the intuitive definition of  $\operatorname{snake}_{i,j}^h$  following (6.4) we see that moving the witness from i-1 to i+1 increases the depth of the witness by one, and similarly moving from i+1 to i-1 decreases the depth by one. As the stars on i and j move in or out along their respective ribbons, one star necessarily moves to an abutting ribbon joined on an east/west edge and the other star moves to abutting ribbon joined on a north/south edge. Moving a star across a north/south edge will not change the depth of the star, but moving a star across an east or west edge will increase or decrease the depth by one, respectively, canceling the contribution from moving the witness. Therefore the total contribution to spin from i-1, i, i+1 remains the same after applying snake $_{i,j}^h$ . All cases are now covered.

The results in Theorem 6.4 and Proposition 6.5 naturally extend to skew partitions as well since the proofs only involve intervals of rank 3 in the *n*-core poset.

Corollary 6.6. Let  $\mu \subset \nu$  be n-cores of lengths  $\ell(\mu) = p$  and  $\ell(\nu) = q$ . Then, for p < i < q, the map  $\varphi_i$  is a well-defined, spin preserving, involution on all starred strong skew tableaux for  $\nu/\mu$ . In particular, spin is constant on affine dual equivalence classes.

#### 7. A Graph on Starred Strong Tableaux

In this section, we construct a vertex-signed, edge-colored graph from our elementary affine dual equivalence map  $\varphi_i$ . The main goal of this section is to show that this graph is, in fact, an LSP<sub>2</sub> graph by Definition 4.5. In order to establish this, we introduce two operations on starred strong tableaux which together show that there are only finitely many isomorphism types for 2-colored connected components. The reduction to finitely many isomorphism types can be viewed as an (incomplete) analog of the jeu da taquin algorithm for starred strong tableaux. This analogy is summarized in Remark 7.12.

**Definition 7.1.** For an *n*-core  $\nu$ , the affine dual equivalence graph  $\mathcal{G}_{\nu}^{(n)}$  is the signed, colored graph with vertex set  $V_{\nu}$  given by the set of all starred strong tableaux  $S^*$  of

shape  $\nu$ , with signature function  $\sigma(S^*)$  obtained from the reading word on the starred letters in  $S^*$ , and for each  $1 < i < \ell(\nu)$ , the set of *i*-colored edges,  $E_i$ , is the set of all pairs  $\{S^*, \varphi_i(S^*)\}$  such that  $S^* \neq \varphi_i(S^*)$ . This definition also extends to skew shapes  $\nu/\mu$  in the *n*-core poset. For  $S^* \in \text{SST}^*(\nu/\mu, n)$ , let  $[S^*]$  denote the connected component of the affine dual equivalence graph  $\mathcal{G}_{\nu/\mu}^{(n)}$  containing  $S^*$ .

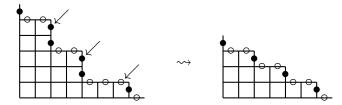
For example, for n = 3 and  $\mu = (5, 3, 1)$  the affine dual equivalence graph is shown on page 47.

Recall that  $\varphi_i$  is an involution which preserves the spin statistic by Corollary 6.6. In order to justify our terminology of affine dual equivalence. We want to prove that the graph induced by these involutions satisfies Axioms 1,2,3,5 from Definition 4.2 and local Schur positivity on all two adjacent colored connected components. Thus each affine dual equivalence graph is LSP<sub>2</sub>. The key will be reducing local Schur positivity to a finite verification. The reduction is achieved with the help of flattening rows and squashing and/or cloning columns.

7.1. The flattening map. Here we define an iterative procedure to flatten an n-core partition down to an m-core partition for any  $1 \le m < n$ . We will extend this procedure to starred strong tableaux in a way that commutes with the affine dual equivalence involutions.

**Definition 7.2.** For any m+1-core  $\lambda$  and any  $1 \leq d \leq m+1$ , define  $\lambda^{(d)}$  to be the unique partition associated to the binary string obtained by removing all beads and spacers with content congruent to d modulo m+1 from the abacus of  $\lambda$ . In particular,  $\lambda^{(d)}$  is an m-core.

We note that the above definition makes sense in light of Remark 5.1 and the characterization of *n*-cores in terms of the *n*-rod abacus. For example, regarding (7,4,4,2,2) as a 4-core,  $(7,4,4,2,2)^{(2)}$  is the 3-core (6,4,2).



Remark 7.3. For n-cores  $\mu \subset \nu$ , if some transposition sequence from  $\mu$  to  $\nu$  touches rod d then every transposition sequence from  $\mu$  to  $\nu$  touches rod d. This follows from the observation that any saturated chain from  $\mu$  to  $\nu$  can be obtained from any other by some sequence of interval exchanges, none of which may change which rods are touched.

**Proposition 7.4.** Let  $\mu \subset \nu$  be m+1-cores such that some (equivalently, every) transposition sequence from  $\mu$  to  $\nu$  does not touch rod d. Then the interval  $[\mu, \nu]$  in the m+1-core poset is isomorphic to  $[\mu^{(d)}, \nu^{(d)}]$  in the m-core poset. This isomorphism extends to a bijection on skew strong tableaux which preserves the number of i-ribbons for each i.

Remark 7.5. Proposition 7.4 can be used in reverse: given  $\mu < \nu$  in the m-core poset, we can lift the interval  $[\mu, \nu]$  to an isomorphic interval in the m+1-core posets with the same nice implications on strong tableaux. This map is implemented by using the inverse procedure of adding in an extra rod between any two existing rods. This can be done precisely when the length of the inserted rod never has length weakly between the length of two interchanging rods, for instance, we may always take the rod to be longer than all other rods or shorter than all other rods.

*Proof.* Recall that exchanging rods in the n-rod abacus preserves the fact that the corresponding binary strings are balanced. Since the length of rod d for each m+1-core  $\lambda$  in the interval  $[\mu, \nu]$  is constant, the re-indexing for each  $\lambda^{(d)}$  is the same. Further, since the covering relations in the m+1-core poset depend on rod d only in the sense that it must not have length weakly between that of the two exchanging rods, covering relations in mapping  $[\mu, \nu]$  down to the m-core poset are preserved. Conversely, given any m-core  $\gamma \in [\mu^{(d)}, \nu^{(d)}]$ , we can lift it to an m+1-core by reversing the procedure. The reverse procedure also is injective and preserves containment order. Hence the intervals are isomorphic.

The bijection on skew strong tableaux is obtained in the obvious way, by mapping the saturated chain

$$S = (\mu = \mu_0 \subset \mu_1 \subset \mu_2 \subset \cdots \subset \mu_k = \nu)$$

to the chain

$$S^{(d)} = (\mu^{(d)} = \mu_0^{(d)} \subset \mu_1^{(d)} \subset \mu_2^{(d)} \subset \dots \subset \mu_k^{(d)} = \nu^{(d)}).$$

To see that this bijection preserves the number of *i*-ribbons, recall from Corollary 5.4 that the number of *i*-ribbons of a strong tableau is equal to the difference in length of the interchanging rods taking  $\lambda_{i-1}$  to  $\lambda_i$ . Since the map from m+1-cores to m-cores preserves the relative lengths of all rods other than rod d, this number is clearly preserved.

By Proposition 7.4, the following map is well defined.

**Definition 7.6.** Let  $\mu \subset \nu$  be m+1-cores such that some (equivalently, every) transposition sequence from  $\mu$  to  $\nu$  does not touch rod d. Define the *flattening map* 

$$fl_d: SST^*(\nu/\mu, m+1) \longrightarrow SST^*(\nu^{(d)}/\mu^{(d)}, m)$$

sending  $S^* \in SST^*(\nu/\mu, m+1)$  to the underlying strong tableau  $S^{(d)}$  with the stars placed on each *i*-ribbon in such a way as to preserve the depth.

Note that the flattening map does not, in general, preserve the spin statistic because it can shorten the height of ribbons.

**Proposition 7.7.** Let  $\mu \subset \nu$  be m+1-cores such that some transposition sequence from  $\mu$  to  $\nu$  does not touch rod d. The flattening map  $\mathfrak{fl}_d : \mathrm{SST}^*(\nu/\mu, m+1) \longrightarrow \mathrm{SST}^*(\nu^{(d)}/\mu^{(d)}, m)$  is a bijection preserving the signature of a starred strong tableau and it commutes with the involutions  $\varphi_i$  for all  $1 < i < \ell(\nu) - \ell(\mu)$ .

Proof. To see  $fl_d$  preserves the signature  $\sigma(S^*)$ , recall from Definition 5.5 and Corollary 5.6 that the content of  $i^*$  is determined by an excess bead on the longer rod in the ith exchange on the n-rod abacus. Since the relative order among the beads on the abacus is unchanged by the procedure in Definition 7.2, the contents of  $i^*$ ,  $(i+1)^*$  will form a decent in  $\sigma(S^*)$  if and only if there is a corresponding descent in  $\sigma(fl_d(S^*))$ . This proves  $\sigma(S^*) = \sigma(fl_d(S^*))$ .

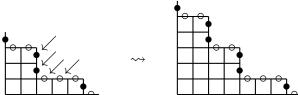
To show  $fl_d(\varphi_i(S^*)) = \varphi_i(fl_d(S^*))$ , simply note that the cases in the definition of  $\varphi_i$  depend only on the types of rod exchanges in the corresponding 3-interval of the m+1 or m-core poset respectively. But, the relative order among the endpoints of the exchanging rods and the isomorphism type of the interval are persevered by the flattening map. Hence the flattening map and the involution commute.

Corollary 7.8. Let  $\mu \subset \nu$  be n-cores with  $\nu$  lying r ranks above  $\mu$ . Then for m=2r, there exists m-cores  $\widehat{\mu} \subset \widehat{\nu}$  such that there exists a bijection from  $\mathrm{SST}^*(\nu/\mu,n)$  to  $\mathrm{SST}^*(\widehat{\nu}/\widehat{\mu},m)$  that preserves the signature and commutes with the involutions  $\varphi_i$  for all  $1 < i < \ell(\nu) - \ell(\mu)$ . Thus, the affine dual equivalence graphs of  $\widehat{\nu}/\widehat{\mu}$  and  $\nu/\mu$  are isomorphic as signed colored graphs.

7.2. The cloning map. Whereas flattening removes rows of the abacus, cloning adds columns. Analogous to flattening, we will define cloning on starred strong tableaux so that it preserves the signatures. In some cases, cloning commutes with the affine dual equivalence operators  $\varphi_i$ .

**Definition 7.9.** For any n-core  $\mu$ , define  $\mu_{(j)}$  to be the unique partition associated to the abacus obtained by *cloning* the column of the n-rod abacus of  $\mu$  containing positions  $j, j+1, \ldots, j+n-1$ . Specifically, let  $\beta$  be the binary string encoding the abacus for  $\mu$ . Then  $\mu_{(j)}$  is the abacus associated to the string obtained from  $\beta$  by inserting a copy of the substring  $\beta_j, \beta_{j+1}, \ldots, \beta_{j+n-1}$  into the abacus for  $\mu$  between positions j-1 and j.

Cloning a column has the effect of extending some of the rods in the n-rod abacus, hence  $\mu_{(j)}$  is also an n-core. To see the effect of cloning on partitions, consider taking (5,2,2) regarded as a 4-core and cloning the column beginning with content 0. This gives  $(5,2,2)_{(0)} = (7,4,4,2,2)$ , as depicted below.



Observe that for fixed  $\mu$ , different values for j can lead to the same n-core  $\mu_{(j)}$ . For instance, taking any  $j \in \{-4, \ldots, 0\}$  results in  $(5, 2, 2)_{(j)} = (7, 4, 4, 2, 2)$ .

In order for flattening to preserve a covering relation in the n-core poset, the transposition sequence simply needs to avoid the rod being removed. The situation for cloning is more subtle. Covering relations are not always preserved even when the replicated column is disjoint from the indexing transposition. It is immediate from Proposition 5.3 that if  $t_{r,s}\mu > \mu$  is a covering relation in the n-core poset, then

 $(t_{r,s}\mu)_{(j)}$  covers  $\mu_{(j)}$  in the *n*-core poset if and and only if for every  $r < h \le s$  the relative order of the lengths of rods h, r and s is the same in both  $\mu$  and  $\mu_{(j)}$ . We call such a j a cloneable index for  $\mu \subset t_{r,s}\mu$ . More generally, j is a cloneable index for the interval  $[\mu, \nu]$  provided cloning the column beginning at j of every core partition in the interval results in another isomorphic interval in the n-core poset. This happens if and only if no rod in the n-rod abacus representing any element in the interval has a rightmost bead of content  $j, j+1, \ldots, j+n-1$ . Similarly, we say j is a cloneable index for  $S^* \in \mathrm{SST}^*(\nu/\mu, n)$  provided j is a cloneable index for  $[\mu, \nu]$ . The clone of  $S^*$ , denoted  $\mathrm{cl}_j(S^*)$ , is defined to be the saturated chain obtained from S by cloning, the column beginning with j in each n-rod abacus in the chain and leaving all the stars with content less than j at the same depth and increasing the depth by 1 for all stars with content at least j. Note, all i-ribbons will have the same shape in  $S^*$  and  $\mathrm{cl}_j(S^*)$  since the relative order of the rod lengths is unchanged by cloning a column. See Figure 11 for example.

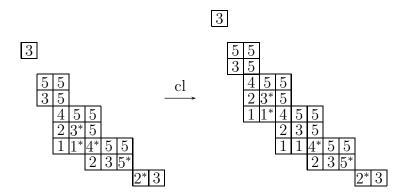


FIGURE 11. An example of the cloning map on a starred strong tableau.

Observe that if  $S^* \in SST^*(\nu/\mu, n)$  and j is a cloneable index for  $S^*$ , then j is a cloneable index for every other starred strong tableau in  $SST^*(\nu/\mu, n)$  as well since the definition of a cloneable index only depends on the interval  $[\mu, \nu]$ .

**Definition 7.10.** Assume that  $S^*$  has a cloneable index at j and that  $T^* = \operatorname{cl}_j(S^*) \in \operatorname{SST}^*(\beta/\alpha, n)$ . Define the *cloning map* on components

$$\operatorname{cl}_i: [S^*] \longrightarrow \operatorname{SST}^*(\beta/\alpha, n)$$

by cloning each starred strong tableaux in  $[S^*]$  at the column beginning with j. The inverse map to cloning, when its defined, will be denoted by

$$\operatorname{sq}_i: [T^*] \longrightarrow [S^*]$$

and we call it the squashing map.

As with the flattening map, the cloning map does not, in general, preserve the spin statistic since it may alter the number of *i*-ribbons and/or it may alter the depth of the starred ribbons. Nonetheless, once the cloning map commutes with the  $\varphi_i$ 's on a connected component of an affine dual equivalence graph then we can clone the same column any number of times and get an isomorphic component.

The following proposition is the analog of Proposition 7.7.

**Proposition 7.11.** Assume that  $S^* \in SST^*(\nu/\mu, n)$  has a cloneable index at j and that  $T^* = \operatorname{cl}_j(S^*)$ . Further assume that cloning the column beginning at j commutes with the involutions  $\varphi_i$  for all  $1 < i < \ell(\nu) - \ell(\mu)$  on the component  $[S^*]$ . Then j is a cloneable index for every starred strong tableaux in  $[T^*]$ . Moreover, if  $U^* = \operatorname{cl}_j(T^*)$ , then  $\operatorname{cl}_j : [T^*] \longrightarrow [U^*]$  is a bijection preserving the signature of each starred strong tableau and it commutes with the involutions  $\varphi_i$  for all  $1 < i < \ell(\nu) - \ell(\mu)$ . Thus,  $[S^*] \approx [T^*] \approx [U^*]$  as signed, colored graphs.

*Proof.* The fact that j is again a cloneable index for  $T^*$  follows directly from the characterization of j being a cloneable index for the interval containing  $S^*$  in terms of rod lengths.

To see that  $[T^*]$  is isomorphic to  $[U^*]$  as signed colored graphs, one must check that the affine dual equivalence maps  $\varphi_i$  commute with the cloning map from  $[T^*]$  to  $[U^*]$ . This follows since the conditions for the affine dual equivalence map on rank 3 intervals are unchanged at each step by removing n consecutive content diagonals that contains no head or tail of a starred ribbon in any of the starred strong tableaux in  $[U^*]$  and that the cells in those n diagonals must necessarily be a copy of the next n-translate down.

- Remark 7.12. As a consequence of Proposition 7.7 and Proposition 7.11, we observe that the process of flattening and squashing a component in an affine dual equivalence graph as much as possible is similar to applying the necessary jeu da taquin slides which bring together all of the connected components in a skew tableaux by removing empty rows and columns. Note, both flattening and cloning/squashing can change the spin statistic even when they commute with affine dual equivalence on a component. Thus a complete analog of jeu da taquin generalizing these moves would need to keep track of powers of t separately from the algorithm.
- 7.3. **Local Schur positivity.** Our next goal is to show that there are only a small number of isomorphism classes of connected components of rank 4 affine dual equivalence graphs. Recall that a starred strong tableau  $S^*$  on an interval  $[\mu, \nu]$  has ribbons labeled 1,2, ...,  $\ell(\nu) \ell(\mu)$ . We say  $S^*$  has  $rank\ r$  provided  $r = \ell(\nu) \ell(\mu)$ . The component  $[S^*]$  of the affine dual equivalence graph on  $\nu/\mu$  has edges labeled 2,3,.., r-1 and each vertex has a signature of length r-1.
- **Lemma 7.13.** Let  $S^* \in SST^*_n(\nu/\mu)$  be a starred strong tableau of rank k = 4. Then  $[S^*]$  has a Schur positive generating function. In fact, each such  $[S^*]$  is either an isolated vertex, or a path with either 2 or 4 edges with alternating color labels. See Figure 12.

This lemma can be proved in two ways. One approach is to do a computer verification by identifying a set of dual equivalence classes which contain all possible isomorphism types after flattening and squashing as much as possible. Details of this approach can be found at http://www.math.washington.edu/~billey/kschur/. The second approach is based on the reading words of the starred strong tableaux, see [3].

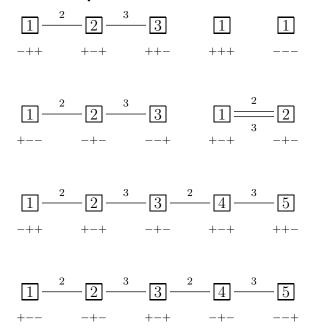


FIGURE 12. All 7 possible isomorphism types of connected components of affine dual equivalence graphs of rank 4.

Remark 7.14. A computer exploration for all possible isomorphism types for affine dual equivalence graphs of rank 5 is underway. As of November of 2011, we have observed 326 distinct isomorphism types which can be viewed in http://www.math.washington.edu/~billey Note for comparison, there are only 25 isomorphism types for rank 5 graphs for LLT polynomials as defined in Section 8.

**Theorem 7.15.** For any pair of n-core partitions  $\mu \subset \nu$ , the affine dual equivalence graph  $\mathcal{G}_{\nu/\mu}^{(n)}$  is a D graph which is locally Schur positive for 2-colored edges and for which spin is constant on connected components.

*Proof.* By Proposition 6.5, the involutions  $\varphi_i$  preserve the spin statistic, hence spin is constant on connected components of  $\mathcal{G}_{\nu/\mu}^{(n)}$ .

To prove  $\mathcal{G}_{\nu/\mu}^{(n)}$  is a D graph, we must verify the axioms in Definition 4.5. Axiom 1 follows from Theorem 6.4 where  $\varphi_i$  is shown to be an involution which switches the sign appropriately. Axioms 2 and 5 follow from the fact that  $\varphi_i$  affects only i-1,i and i+1-ribbons. Axiom 3 and the LSP<sub>2</sub> property both follow from Lemma 7.13 since every connected component of  $\mathcal{G}_{\nu/\mu}^{(n)}$  restricted to  $E_{i-1} \cup E_i$  is isomorphic to a component of a rank 4 affine dual equivalence graph replacing the edge labels 2,3 by i-1,i respectively.

Note that affine dual equivalence graphs need not satisfy Axiom 4 of Definition 4.2. It is not known if affine dual equivalence graphs satisfy Axiom 6.

## 8. Connections with LLT and MacDonald Polynomials

The primary interest in k-Schur functions originally was the conjecture of Lapointe, Lascoux and Morse that these functions straddle the gap between Macdonald polynomials and Schur functions. That is, when  $\mu$  is a k-bounded partition, they conjecture

(8.1) 
$$H_{\mu}(X;q,t) = \sum_{\nu < \mu} K_{\nu,\mu}^{(k)}(q,t) \ s_{\nu}^{(k)}(X;t),$$

where  $K_{\nu,\mu}^{(k)}(q,t) \in \mathbb{N}[q,t]$ , and

(8.2) 
$$s_{\nu}^{(k)}(X;t) = \sum_{\lambda < \nu} C_{\lambda,\nu}^{(k)}(t) \ s_{\lambda}(X),$$

where  $C_{\lambda,\nu}^{(k)}(t) \in \mathbb{N}[t]$ .

Using the definition of k-Schur functions advocated for in this paper, we show how our methods shed light on equation (8.1), and, more generally, the problem of expanding LLT polynomials into k-Schur functions.

8.1. Macdonald polynomials. The transformed Macdonald polynomials  $\widetilde{H}_{\mu}(X;q,t)$  form a basis for symmetric functions with two additional parameters. Precisely,  $\{\widetilde{H}_{\mu}(X;q,t)\}$  is a basis for  $\Lambda$  with coefficients in  $\mathbb{Q}(q,t)$ . Macdonald [27] originally defined the polynomials to be the unique functions satisfying certain orthogonality and triangularity conditions. Haglund's monomial (quasisymmetric) expansion for Macdonald polynomials [9, 10] gives an explicit combinatorial description of  $\widetilde{H}_{\mu}(X;q,t)$  as the q,t-generating function of permutations, regarded as standard fillings of the diagram of  $\mu$ .

For a cell x in the diagram of a partition  $\mu$ , let l(x) (respectively a(x)) denote the number of cells directly north (respectively east) of x. Given a permutation w of  $\{1, 2, \ldots, |\mu|\}$ , fill the diagram of  $\mu$  with w written in one-line notation so that w becomes the row reading word of the resulting filling. A  $\mu$ -descent of such a filling is a pair of cells (x, y) with x immediately north of y and the entry in x is greater than the entry in y. Denote by  $\operatorname{Des}_{\mu}(w)$  the set of all  $\mu$ -descents of w. Define the major index with respect to  $\mu$  to be

(8.3) 
$$\mathrm{maj}_{\mu}(w) = \sum_{(x,y) \in \mathrm{Des}_{\mu}(w)} l(x) + 1.$$

Note that when  $\mu$  is a single column, maj $_{\mu}$  is the usual major index on permutations. An ordered pair of cells (x,y) in the diagram of  $\mu$  is called *attacking* if x and y lie in the same row with x strictly west of y, or if x is in the row immediately north of y and x lies strictly east of y. Given a permutation filling of  $\mu$ , a  $\mu$ -inversion pair is an attacking pair (x,y) where the entry of x is greater than the entry of y. Denote by  $\text{Inv}_{\mu}(w)$  the set of inversion pairs of w filled into  $\mu$ ; this set is a subset of the usual inversion set for w. Define the inversion number with respect to  $\mu$  to be

(8.4) 
$$\operatorname{inv}_{\mu}(w) = |\operatorname{Inv}_{\mu}(w)| - \sum_{(x,y) \in \operatorname{Des}_{\mu}(w)} a(x).$$

Note that when  $\mu$  is a single row, inv<sub> $\mu$ </sub> is the usual inversion number on permutations.

For example, let  $\mu$  be the partition (5, 4, 4, 1) and take  $w = [5 \ 11 \ 14 \ 9 \ 2 \ 6 \ 3 \ 4 \ 10 \ 8 \ 1 \ 13 \ 7 \ 12]$ in  $S_{14}$ . Filling w into  $\mu$  gives

Abusing notation, represent a cell of the filling by the entry which it contains. The  $\mu$ -descent set of w is

$$Des_{\mu}(w) = \{(11,6), (14,3), (3,1), (9,4), (10,7)\},\$$

and the  $\mu$ -inversion pairs of w are given by

$$\operatorname{Inv}_{\mu}(w) = \left\{ \begin{array}{lll} (11,9), & (14,2), & (9,6), & (6,4), & (10,1), & (13,7), \\ (11,2), & (14,6), & (9,3), & (4,1), & (8,1), & (13,12) \\ (14,9), & (9,2), & (6,3), & (10,8), & (8,7), \end{array} \right\}.$$

Therefore the maj<sub> $\mu$ </sub> and inv<sub> $\mu$ </sub> statistics associated to w are

$$\operatorname{maj}_{\mu}(w) = 2 + 1 + 2 + 1 + 2 = 8,$$

$$\operatorname{inv}_{\mu}(w) = 17 - (3 + 2 + 2 + 1 + 0) = 9.$$

Remark 8.1. If  $(x,y) \in Des_{\mu}(w)$ , then for every cell z of the arm of x, the entry of z is either bigger than the entry of y or smaller than the entry of x (or both). In the former case, (z, y) will form an inversion pair, and in the latter case, (x, z) will form an inversion pair. Therefore  $\operatorname{inv}_{\mu}(w)$  is a non-negative integer.

Define the signature function 
$$\sigma: \mathcal{S}_n \to \{\pm 1\}^{n-1}$$
 on permutations by
$$\sigma_i(w) = \begin{cases} +1 & \text{if } i \text{ lies left of } i+1 \text{ in } w \\ -1 & \text{if } i+1 \text{ lies left of } i \text{ in } w. \end{cases}$$

For the permutation above,  $\sigma(w) = -++-++--+-$  if we abbreviate  $\{-1, +1\}$ by  $\{-,+\}$ . Using  $\sigma$  to associate a quasisymmetric function to each permutation filling of  $\mu$ , Haglund's formula for  $\widetilde{H}_{\mu}(X;q,t)$  may be stated as follows.

**Definition 8.2.** [9] The transformed Macdonald polynomials are given by

(8.7) 
$$\widetilde{H}_{\mu}(X;q,t) = \sum_{w \in \mathcal{S}_n} q^{\operatorname{inv}_{\mu}(w)} t^{\operatorname{maj}_{\mu}(w)} Q_{\sigma(w)}(X).$$

It is a theorem in [10] that (8.7) satisfies the conditions which uniquely characterize the transformed Macdonald polynomials as originally defined in [27]. The proof is by an elegant and elementary combinatorial argument, so we take Haglund's formula as the definition.

A combinatorial proof of Macdonald positivity is given in [1] by putting a D graph structure on permutation fillings of a partition diagram. In this case, the edges of the graph are defined by simple involutions on the permutations. The *i-witness* is defined as usual to be the middle letter of i-1, i, i+1 encountered when reading the permutation from left to right. The main ingredients in the edges are usual dual equivalence on permutations, denoted  $d_i$ , and a natural modification of dual equivalence, denoted  $\widetilde{d_i}$ . Precisely, these involutions are defined by

(8.8) 
$$\cdots i \cdots i \pm 1 \cdots i \mp 1 \cdots \stackrel{d_i}{\longleftrightarrow} \cdots i \mp 1 \cdots i \pm 1 \cdots i \cdots,$$

(8.9) 
$$\cdots i \cdots i \pm 1 \cdots i \mp 1 \cdots \stackrel{\widetilde{d_i}}{\longleftrightarrow} \cdots i \pm 1 \cdots i \mp 1 \cdots i$$

Note that in the former case the *i*-witness always remains the same while in the latter it always toggles between i-1 and i+1. The edge-defining involutions of the D graph for  $\widetilde{H}_{\mu}(X;q,t)$  is then given by

(8.10) 
$$\varphi_{i}^{\mu}(w) = \begin{cases} w & \text{if } i \text{ is the } i\text{-witness,} \\ \widetilde{d}_{i}(w) & \text{if } i - 1, i, i + 1 \text{ fit in } \\ d_{i}(w) & \text{otherwise.} \end{cases}$$

A key observation in [1] is that  $\varphi_i^{\mu}$  preserves Haglund's statistics maj<sub> $\mu$ </sub> and inv<sub> $\mu$ </sub>. Much like the case for starred strong tableaux, the proof that the resulting graph satisfies the Axioms 1,2,3, and 5 is rather straightforward. Axiom 4' requires a simpler computer verification.

8.2. **LLT Polynomials.** We may also regard Haglund's formula for Macdonald polynomials as a weighted sum over tableaux-like objects. In this paradigm, equation (8.7) can be interpreted as giving a positive expansion of  $\tilde{H}_{\mu}(X;q,t)$  in terms of certain *LLT polynomials*. Lascoux, Leclerc and Thibon [26] originally defined  $\text{LLT}_{\lambda}(X;q)$  to be the q-generating function of d-ribbon tableaux of shape  $\mu$  weighted by cospin. Below we give another modified definition that is popular in the literature as the q-generating function of d-tuples of tableaux weighted by d-inversions, first presented in [11]. The equivalence of these definitions uses the abacus model for taking d-cores and d-quotients of partitions [16]; for further details on the correspondence in this context, we refer the reader to [2, 11].

Let  $\lambda$  represent the *d*-tuple of (skew) partitions  $(\lambda^{(0)}, \ldots, \lambda^{(d-1)})$ , each embedded in a specific way in  $\mathbb{N} \times \mathbb{N}$ . For such a *d*-tuple, define the *shifted content* of a cell x by

(8.11) 
$$\widetilde{c}(x) = d \cdot c(x) + i$$

when x is a cell of  $\lambda^{(i)}$ , where c(x) is the usual content of x regarded as a cell of  $\lambda^{(i)}$ . Define the *bandwidth* of a d-tuple to be one plus the difference between the largest and smallest *unshifted* cell contents.

A standard d-tuple of shape  $\lambda$  is a bijective filling of the cells of  $\lambda$  with the letters 1 to m so that entries increase along rows and up columns. For example, below is a standard 5-tuple of shape ((3,2)/(1), (1,1,1), (2,1), (2,2)/(1), (1)), say with the southeasternmost cell of each partition embedded at content 0.

Call a pair of cells (x, y) in a d-tuple  $\lambda$  attacking if  $d > \widetilde{c}(y) - \widetilde{c}(x) > 0$ . For a standard d-tuple  $\mathbf{T}$  of shape  $\lambda$ , define the number of d-inversions, denoted inv $_d(\mathbf{T})$ , to be the number of attacking pairs (x, y) with the entry of x greater than the entry of y. For example, the 5-inversion pairs of the standard 5-tuple above are

$$\left\{ \begin{array}{llll} (11,9), & (14,2), & (9,6), & (6,4), & (10,1), & (13,7), \\ (11,2), & (14,6), & (9,3), & (4,1), & (8,1), & (13,12), \\ (14,9), & (9,2), & (6,3), & (10,8), & (8,7), \end{array} \right\}.$$

Note that if the d-tuple consists of d single boxes, i.e.  $\lambda = ((1), (1), \dots, (1))$ , each embedded to have content 0, then d-inversions are simply the usual inversions in the permutation obtained by reading the entries in increasing order of shifted content.

For a d-tuple  $\lambda$ , define the normalizing constant  $a_{\lambda}$  to be the minimum number of d-inversions of a standard d-tuple of shape  $\lambda$ . This normalization is an artifact of q-counting by d-inversions rather than cospin; see [11].

Remark 8.3. Define an inversion triple to be a triple of cells (x, y, z) such that x lies immediately north of z and y has shifted content between that of x and z. Then both (x, y) and (y, z) are attacking pairs. Since x is north of z, we must have x > z, and so at least one of (x, y) and (y, z) will be a d-inversion. Say that two inversions triples are overlapping if they have the form (w, x, y) and (x, y, z). Note that two overlapping inversion triples of this form may contribute only one d-inversion if w < x, x > y and y < z. Therefore the normalizing constant  $a_{\lambda}$  is also equal to the maximum number of pairwise nonoverlapping inversion triples of  $\lambda$ .

**Definition 8.4.** The *LLT polynomial* of shape  $\lambda$ , denoted  $LLT_{\lambda}$ , is defined by

(8.13) 
$$LLT_{\lambda}(X;q) = \sum_{\mathbf{T} \in SYT_{d}(\lambda)} q^{inv_{d}(\mathbf{T}) - a_{\lambda}} Q_{\sigma(\mathbf{T})}(X),$$

where the sum is over standard d-tuples of shape  $\lambda$ ,  $a_{\lambda}$  is the normalizing constant for  $\lambda$ , and  $\sigma(\mathbf{T})$  is defined analogously to equation (2.5) using shifted contents.

The connection between Macdonald polynomials and LLT polynomials can be seen by transforming the permutation fillings of the diagram of  $\mu$  with a given  $\mu$ -descent set into standard  $\mu_1$ -tuples of a certain shape as follows. Let D be a possible  $\mu$ -descent set. For  $i=1,\ldots,\mu_1$ , let  $\mu_D^{(i-1)}$  be the ribbon obtained from the ith column of  $\mu$  by putting the entry of cell (i,j) immediately south of the entry of cell (i,j+1) if  $((i,j+1),(i,j)) \in D$  and immediately east otherwise. Embed each  $\mu_D^{(i)}$  so that the southeasternmost cell has content 0 and, equivalently, shifted content i. Then each permutation filling of shape  $\mu$  with  $\mathrm{Des}_{\mu} = D$  may be regarded as a standard  $\mu_1$ -tuple of tableaux of shape  $\mu_D$ . For example, the filling of (5,4,4,1) in equation (8.5) corresponds to the standard 5-tuple given in equation (8.12).

Since the major index statistic depends only on the  $\mu$ -descent set, we may define the major index of a descent set D of  $\mu$  by  $\mathrm{maj}_{\mu}(D) = \mathrm{maj}_{\mu}(w)$  for any permutation w with  $\mathrm{Des}_{\mu}(w) = D$ . Similarly, define the arm of a descent set D by  $a(D) = \sum_{(x,y)\in D} a(x)$ . Comparing attacking pairs in both paradigms leads to the following expansion of Macdonald polynomials in terms of LLT polynomials.

**Theorem 8.5.** [10] Macdonald polynomials may be expressed in terms of LLT polynomials as

(8.14) 
$$\widetilde{H}_{\mu}(X;q,t) = \sum_{D} t^{\text{maj}_{\mu}(D)} q^{a_{\mu_{D}} - a(D)} \text{LLT}_{\mu_{D}}(X;q),$$

where the sum is over all possible  $\mu$ -descent sets D.

Note that a(D) counts certain nonoverlapping inversion triples of  $\mu_D$ , hence by Remark 8.3,  $a_{\mu_D} \geq a(D)$ . Therefore equation (8.14) gives a *positive* expansion of Macdonald polynomials in terms of LLT polynomials.

The theory of dual equivalence graphs is used in [2] to establish LLT positivity, and the graph for Macdonald polynomials presented in [1] appears as a special case. The graph for LLT polynomials may be described in terms of the same elementary operations,  $d_i$  and  $\tilde{d}_i$ , on standard d-tuples of tableaux. Define the i-witness of the dual equivalence for i - 1, i, i + 1 to be whichever of i - 1, i, i + 1 has shifted content between the other two. As it transpires, none of the three may have equal shifted contents. The analog of dual equivalence for standard d-tuples is given by

(8.15) 
$$\varphi_i^d(w) = \begin{cases} w & \text{if } i \text{ is the } i\text{-witness,} \\ \widetilde{d}_i(w) & \text{if } |\widetilde{c}(i) - \widetilde{c}(i-1)| \le d \text{ and } |\widetilde{c}(i) - \widetilde{c}(i+1)| \le d, \\ d_i(w) & \text{otherwise.} \end{cases}$$

It is shown in [2] that  $\varphi_i^d$  preserves the number of d-inversions and that the graph constructed from these involutions is in fact a D graph. A close inspection of equations (8.10) and (8.15) reveals that if S is a permutation filling of  $\mu$  and  $\mathbf{T}$  is the standard d-tuple corresponding to S via the bijection described above, then  $\varphi_i^d(\mathbf{T})$  is the standard d-tuple corresponding to the permutation filling  $\varphi_i^{\mu}(S)$  of  $\mu$ .

8.3. Expansions into k-Schur functions. Consider the case when the Macdonald polynomial  $\widetilde{H}_{\mu}(X;q,t)$  is equal to a single LLT polynomial  $\mathrm{LLT}_{\mu_D}$ . This happens when  $\mu$  is a single row, and so  $\mu_D = ((1), \ldots, (1))$  each embedded at content 0. For this extreme case, an LLT polynomial will have the most terms in the quasisymmetric expansion. Similarly, a k-Schur function has the most terms in the quasisymmetric expansion when k is as small as possible, i.e. n = k + 1 = 2. In both cases, the D graph will have no double edges and will always toggle the i-witness across an i-edge.

More to the point, define a map  $\theta$  from starred strong tableaux on the 2-core  $(m, m-1, \ldots, 1)$  to standard tableaux on the m-tuple  $((1), \ldots, (1))$  embedded so that each cell has content 0 as follows. Assuming relative positions for 1 up to i-1 have been chosen, reading from left to right, place i in position  $d(i^*)$ . Once all letters are placed, fill the permutation into the m-tuple. For example,

**Theorem 8.6.** The map  $\theta$  is a D graph isomorphism between the graph of starred strong tableaux on the 2-core  $(m, m-1, \ldots, 1)$  and the graph of standard filling of the m-tuple  $((1), \ldots, (1))$  each embedded at content 0. Furthermore,  $q^{\binom{m}{2}} - \text{spin}(S^*) = \text{inv}(\theta(S^*))$ .

Proof. Since  $d(i^*)$  may be recovered for each i from the permutation,  $\theta$  is clearly a bijection on the underlying vertex sets. Moreover,  $\theta$  will place i to the left of i-1 if and only if  $d(i^*) \leq d(i-1^*)$  which is the case if and only if  $i^*$  lies on an earlier diagonal from  $(i-1)^*$ . Therefore  $\theta$  preserves the relative signatures. A similar analysis of  $d(i-1^*)$  and  $d(i+1^*)$  reveals that the witness for the action on the staircase is precisely the witness for the action on the permutation. Since both actions toggle the witness,  $\theta$  commutes with the respective i-edges of the graphs, and hence is an isomorphism. Finally, adding i at position  $d(i^*)$  creates exactly  $i - (d(i^*) + 1)$  inversions.  $\square$ 

Motivated by Theorem 8.6, define the *cospin* of a starred strong tableau by

(8.16) 
$$\operatorname{cospin}(S^*) = \sum_{i} n(i) \cdot (w(i) - 1) + n(i) - (d(i^*) + 1),$$

where w(i) is the width of an i-ribbon. Define the modified k-Schur functions, denoted  $\widetilde{s}_{\lambda}^{(k)}(X;q)$ , by

(8.17) 
$$\widetilde{s}_{\lambda}^{(k)}(X;q) = \sum_{S^* \in SST^*(\rho(\lambda),n)} q^{\operatorname{cospin}(S^*)} Q_{\sigma(S^*)}(X).$$

Here we have changed to the parameter q in order to highlight connections with LLT polynomials and Macdonald polynomials. Recall from Section 8.1 that a Macdonald polynomial indexed by a single row is precisely an LLT polynomial where each component is a single cell. Therefore we may interpret this isomorphism of D graphs as the following symmetric function identity.

Corollary 8.7. For  $m \geq 1$ , we have

(8.18) 
$$LLT_{(1),\dots,(1)}(X;q) = \widetilde{H}_{(m)}(X;q,t) = \widetilde{s}_{(1^m)}^{(1)}(X;q).$$

Another illuminating case to consider is when an LLT polynomial is equal to a single Schur function. It is easy to see from the definition that this is the case exactly when the indexing tuple consists of a single partition. Similarly, we have the following characterization for k-Schur functions.

**Proposition 8.8.** A k-Schur function is equal to a single Schur function if and only if the indexing partition has bandwidth at most k.

*Proof.* To say  $\lambda$  has bandwidth k is to say that the rim of  $\lambda$  consists of k cells, and thus  $\lambda$  is an n-core, where n=k+1. Moreover, in this case the n-core poset is isomorphic to Young's lattice. Therefore the strong tableaux of shape  $\lambda$  are precisely the standard tableaux of shape  $\lambda$ , and the contribution to spin is nil. This argument is easily reversed.

In both cases, the D graphs will be the standard dual equivalence graph for the indexing partition. On the level of the symmetric functions, we have the following identity.

Corollary 8.9. For  $\lambda$  a partition with bandwidth at most k, we have

$$LLT_{\lambda}(X;q) = s_{\lambda}(X) = \widetilde{s}_{\lambda}^{(k)}(X;q).$$

Corollaries 8.7 and 8.9 support the following, first conjectured by Mark Haiman [12].

**Conjecture 8.10.** Let  $\lambda = (\lambda^{(0)}, \dots, \lambda^{(m-1)})$  be an m-tuple of partitions of bandwidth at most k. Then

(8.19) 
$$\operatorname{LLT}_{\lambda}(X;q) = \sum_{\mu} c_{\lambda,\mu}^{(k)}(q) \ \widetilde{s}_{\mu}^{(k)}(X;q),$$

where  $c_{\boldsymbol{\lambda},\mu}^{(k)}(q) \in \mathbb{N}[q]$ .

Using the expansion of Macdonald polynomials into certain LLT polynomials, this conjecture implies that a Macdonald polynomial indexed by a partition with at most k rows is k-Schur positive. This statement can be reformulated to recover the original conjecture of Lascoux, Lapointe and Morse (8.1) by interchanging q and t and conjugating the indexing partition.

Corollary 8.11. Assuming Conjecture 8.10, if  $\mu$  is a partition with at most k rows, then

(8.20) 
$$\widetilde{H}_{\mu}(X;q,t) = \sum_{\lambda} \widetilde{K}_{\lambda,\mu}^{(k)}(q,t) \ \widetilde{s}_{\lambda}^{(k)}(X;q),$$

where  $\widetilde{K}_{\lambda,\mu}^{(k)}(q,t) \in \mathbb{N}[q,t]$ . In particular, Macdonald polynomials are k-Schur positive.

## Appendix A. Examples

In this appendix we give the quasisymmetric and Schur expansion for the k-Schur function  $s_{(2,2,1)}^{(2)}$ . We compute this using the interval  $[\emptyset, (5,3,1)]$  of the 3-core poset (Figure 13) and the corresponding D graph on all starred strong tableaux of shape (5,3,1) regarded as a 3-core (Figure 14).

$$\begin{split} s_{(2,2,1)}^{(2)} &= Q_{-+--} \, + \, Q_{--+-} \, + \, t \, Q_{++--} \, + \, (1+t) \, Q_{+--+} \, + \, (2t+t^2) \, Q_{-++-} \\ &\quad + \, (1+2t+t^2) \, Q_{+-+-} \, + \, (1+2t+t^2) \, Q_{-+-+} \, + \, t \, Q_{--++} \\ &\quad + \, (t^2+t^3) \, Q_{+++-} \, + \, (t+2t^2+t^3) \, Q_{++-+} \, + \, (t+2t^2+t^3) \, Q_{+-++} \\ &\quad + \, (t^2+t^3) \, Q_{-+++} \, + \, t^4 Q_{++++} \\ &= s_{(2,2,1)} + t s_{(3,1,1)} + (t+t^2) s_{(3,2)} + (t^2+t^3) s_{(4,1)} + t^4 s_{(5)} \end{split}$$

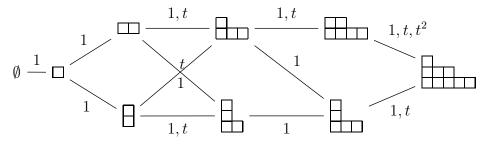


FIGURE 13. The poset of 3-cores lying below (5,3,1), with edge weights giving the spin contributions of possible starrings.

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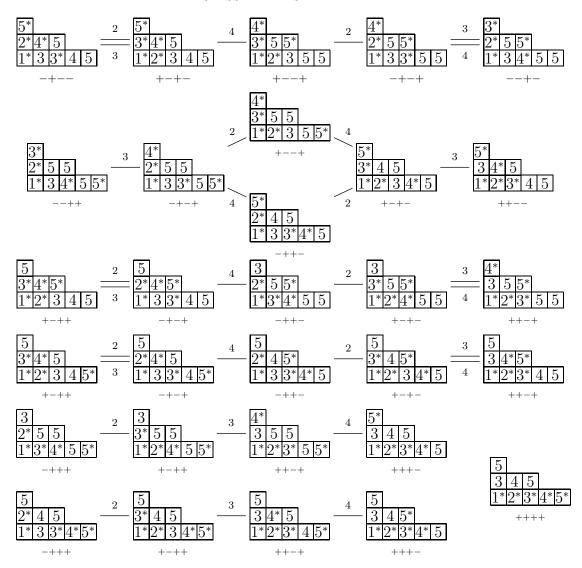


FIGURE 14. The D graph on starred strong tableaux of shape (5,3,1) regarded as a 3-core.

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Department of Mathematics, Massachusetts Institute of Technology, 77 Massachusetts Avenue, Cambridge, MA 02139-4307

E-mail address: sassaf@math.mit.edu

Department of Mathematics, Padelford C-445, University of Washington, Box 354350, Seattle, WA 98195-4350

E-mail address: billey@math.washington.edu